

Cargèse Fall School on Random Graphs
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INTRODUCTION TO RANDOM GRAPHS

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TWO MAIN RANDOM GRAPH MODELS

THE BINOMIAL RANDOM GRAPH $G(n, p)$

$G(n, p)$ is the (random) graph on vertices $\{1, 2, \dots, n\}$ in which each of $\binom{n}{2}$ possible pairs appears as an edge independently with probability p .

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WHAT DOES IT MEAN?

Given a graph G with vertex set $[n]$:

$$\Pr(G(n, p) = G) = p^{e(G)}(1 - p)^{\binom{n}{2} - e(G)}.$$

while

$$\Pr(G(n, M) = G) = \begin{cases} 0 & \text{if } e(G) \neq M \\ 1 / \binom{\binom{n}{2}}{M} & \text{if } e(G) = M \end{cases}$$

ASYMPTOTICS

Typically, we are interested only in the asymptotic behaviour of $G(n, M)$ for very large n , where $M = M(n)$.

For a given function $M = M(n)$, we say that a property holds for $G(n, M)$ **aas** if the probability that it holds for $G(n, M)$ tends to 1 as $n \rightarrow \infty$.

Of course, it is an abuse of language, as in many cases in terminology in the theory of random structures. In fact during this talk I will not be too meticulous in, say, referring to some results – let me apologize for it in advance.

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(Most of) asymptotic properties of $G(n, M)$ and $G(n, p)$ are very similar, provided $p = M/\binom{n}{2}$.

OBSERVATION

Results on $G(n, M)$ are, in a way, more precise, since

$$\Pr(G(n, M) = G) = \Pr(G(n, p) = G | e(G(n, p)) = M),$$

i.e., roughly speaking,

$$G(n, M) = G(n, p) | \{e(G(n, p)) = M\}.$$

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LET US START WITH SOMETHING EASY

THEOREM ERDŐS, RÉNYI'59

Let $p(n) = \frac{1}{n}(\ln n + \gamma(n))$. Then

$$\lim_{n \rightarrow \infty} \Pr(G(n, p) \text{ is connected}) = \begin{cases} 0 & \text{if } \gamma(n) \rightarrow -\infty, \\ 1 & \text{if } \gamma(n) \rightarrow \infty. \end{cases}$$

... OR SOMETHING EVEN EASIER

THEOREM ERDŐS, RÉNYI'59

Let $p(n) = \frac{1}{n}(\ln n + \gamma(n))$. Then

$$\lim_{n \rightarrow \infty} \Pr(\delta(G(n, p)) > 0) = \begin{cases} 0 & \text{if } \gamma(n) \rightarrow -\infty, \\ 1 & \text{if } \gamma(n) \rightarrow \infty. \end{cases}$$

THE FIRST MOMENT METHOD

MARKOV INEQUALITY

Let X be a **non-negative, integer-valued** random variable.

Then

$$\Pr(X > 0) = \Pr(X \geq 1) \leq \mathbb{E}X .$$

THE FIRST MOMENT METHOD

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Let $p(n) = \frac{1}{n}(\ln n + \gamma(n))$ and $\gamma(n) \rightarrow \infty$.

Moreover, let X count isolated vertices in $G(n, p)$.

Then $\Pr(X > 0) \rightarrow 0$ as $n \rightarrow \infty$.

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Proof Note that $X = \sum_{i=1}^n X_i$, where

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In our case

$$\begin{aligned} \mathbb{E}X_i &= (1 - p)^{n-1} = \exp(- (n - 1) \log(1 - p)) \\ &= \exp(-np + O(p + p^2 n)). \end{aligned}$$

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If $p(n) = \frac{1}{n}(\ln n + \gamma(n))$, then

$$\begin{aligned}\mathbb{E}X &= \sum_{i=1}^n \mathbb{E}X_i = n \exp(-np + O(p + p^2n)) \\ &= (1 + o(1))e^{-\gamma},\end{aligned}$$

and so, for $\gamma(n) \rightarrow \infty$, we get

$$\Pr(X > 0) \leq \mathbb{E}X \rightarrow 0. \quad \square$$

REMARK

If we apply the first moment method to the random variable Y which counts non-trivial components in $G(n, p)$ we get a much stronger result.

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If $p(n) = \frac{1}{n}(\ln n + \gamma(n))$, where $\gamma(n) \rightarrow \infty$, then $G(n, p)$ is a.s. connected.

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Quite often (but by no means always) it is the case!

THE SECOND MOMENT METHOD

OBSERVATION

If X counts structures which are “mostly weakly-dependent”, then the expected number of ordered pairs of such structures is roughly $(\mathbb{E}X)^2$, i.e.

$$\mathbb{E}X(X - 1) = (1 + o(1))(\mathbb{E}X)^2.$$

Then, for the variance of X , we have

$$\text{Var}X = \mathbb{E}X(X - 1) + \mathbb{E}X - (\mathbb{E}X)^2 = o(\mathbb{E}X)^2.$$

CHEBYSHEV'S AND CAUCHY'S INEQUALITIES

Let us assume that $\mathbb{E}X \rightarrow \infty$, $\mathbb{E}X(X-1) = (1 + o(1))(\mathbb{E}X)^2$, and so $\text{Var}X = o(\mathbb{E}X)^2$.

CHEBYSHEV'S INEQUALITY

$$\Pr(X = 0) \leq \Pr(|X - \mathbb{E}X| \leq \mathbb{E}X) \leq \frac{\text{Var}X}{(\mathbb{E}X)^2} \rightarrow 0.$$

CAUCHY'S INEQUALITY

If X is an integer-valued, non-negative random variable, then

$$\Pr(X > 0) = \Pr(X \geq 1) \geq \frac{(\mathbb{E}X)^2}{\mathbb{E}X^2} = \frac{(\mathbb{E}X)^2}{\mathbb{E}X(X-1) + \mathbb{E}X} \rightarrow 1.$$

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CHEBYSHEV'S VS. CAUCHY'S

The left hand side of Chebyshev's inequality can be larger than one while Cauchy's bound is always strictly positive!

REVENONS À NOS MOUTONS

Let X be the number of isolated vertices in $G(n, p)$, where $p(n) = \frac{1}{n}(\log n + \gamma(n))$ and $\gamma \rightarrow -\infty$.

Then $\mathbb{E}X = (1 + o(1))e^{-\gamma} \rightarrow \infty$. What about $\mathbb{E}X(X - 1)$?

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$$\begin{aligned}\mathbb{E}X(X - 1) &= n(n - 1)(1 - p)^{2(n-1)-1} = \frac{n - 1}{n(1 - p)} [n(1 - p)^{n-1}]^2 \\ &= (1 + o(1))(\mathbb{E}X)^2.\end{aligned}$$

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Thus, $\Pr(X > 0) \rightarrow 1$. □

ERDŐS-RÉNYI THEOREM IS FINALLY SHOWN!

THEOREM ERDŐS, RÉNYI '59

Let $p(n) = \frac{1}{n}(\ln n + \gamma(n))$. Then

- (I) If $\gamma \rightarrow -\infty$, then aas $G(n, p)$ contains isolated vertices (and so aas it is not connected);
- (II) If $\gamma \rightarrow \infty$, then aas $G(n, p)$ is connected (and so contains no isolated vertices).

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Can we define (and prove) even stronger result which relates connectivity to the absence of isolated vertices?

THE HITTING TIME

THE RANDOM GRAPH PROCESS

$G(n, M)$ can be viewed as the $(M + 1)$ th stage of a Markov chain $\{G(n, M) : 0 \leq M \leq \binom{n}{2}\}$, where we add edges to a graph in a random order.

THE HITTING TIME

Let $h_1 = \min\{M : \delta(G(n, M)) \geq 1\}$ and $h_{\text{conn}} = \min\{M : G(n, M) \text{ is connected}\}$.
Note that both h_1 and h_{conn} are random variables!

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As $h_1 = h_{\text{conn}}$.

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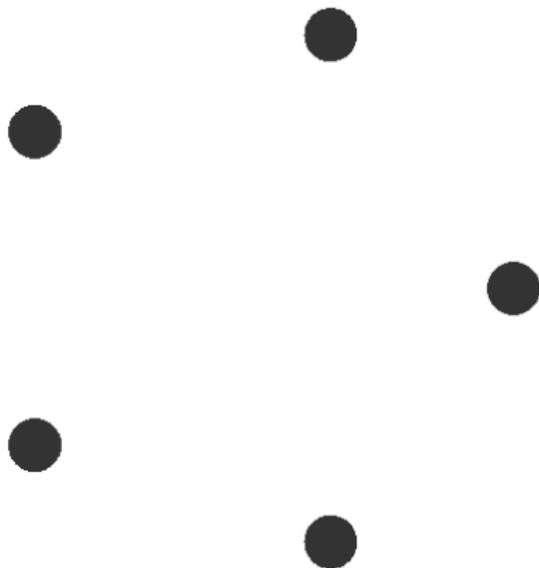
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$$\{G(n, p) : 0 \leq p \leq 1\}$$

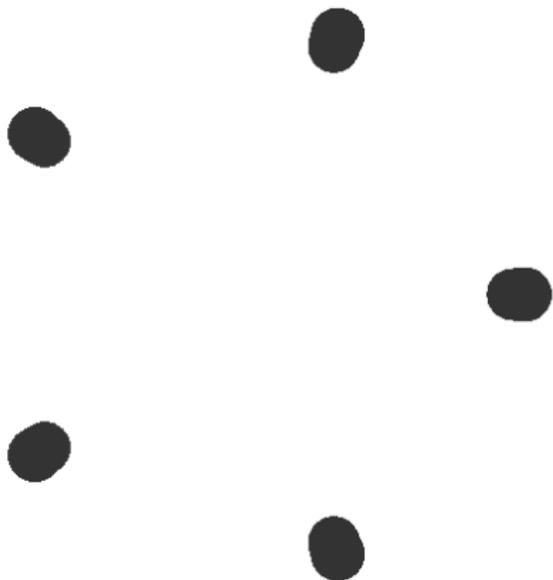
THE RANDOM GRAPH PROCESS (FOR $G(n, p)$)

$G(n, p)$ can also be viewed as a stage of a Markov process $\{G(n, M) : 0 \leq p \leq 1\}$.

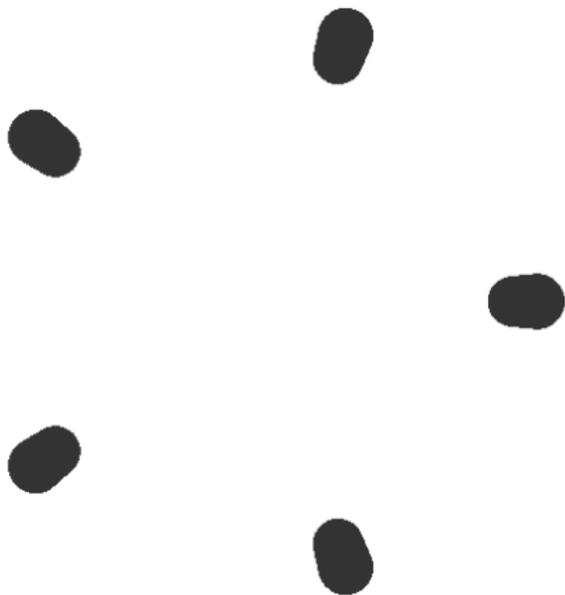
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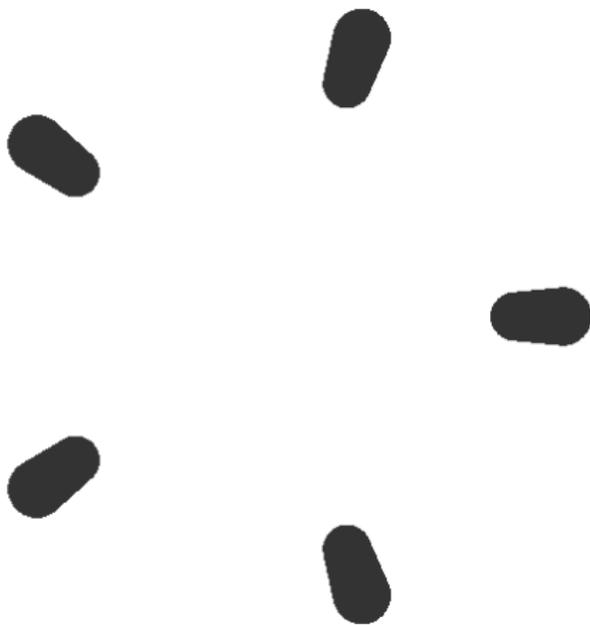
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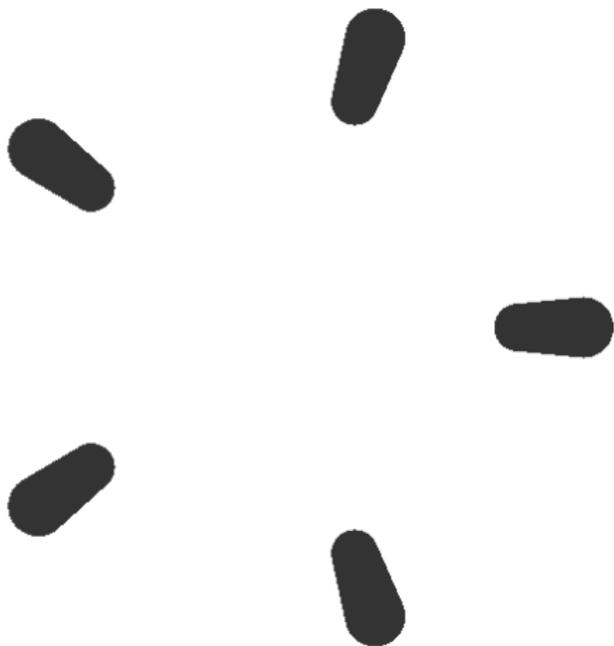
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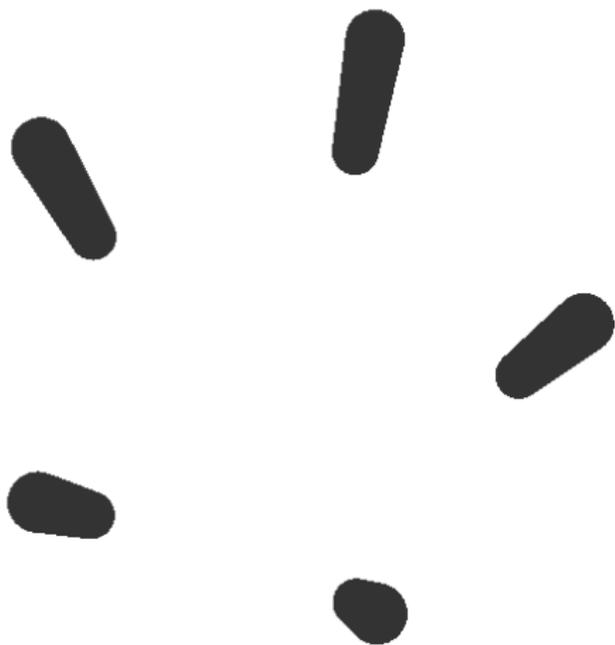
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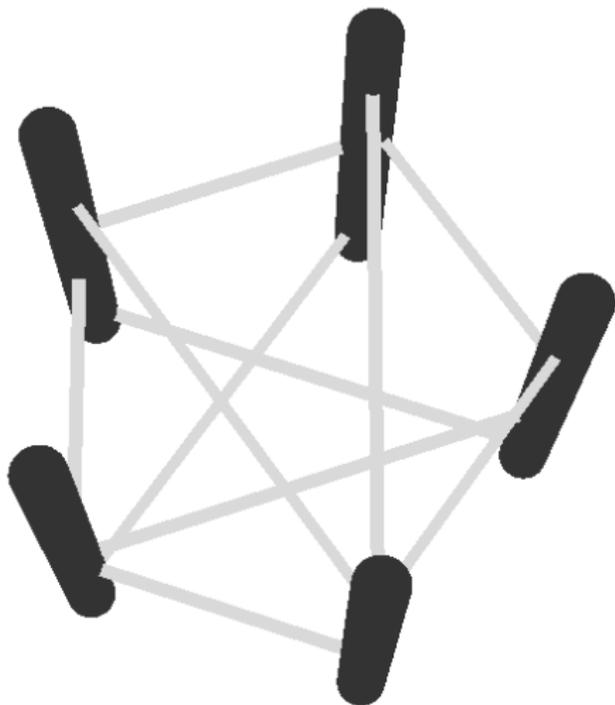
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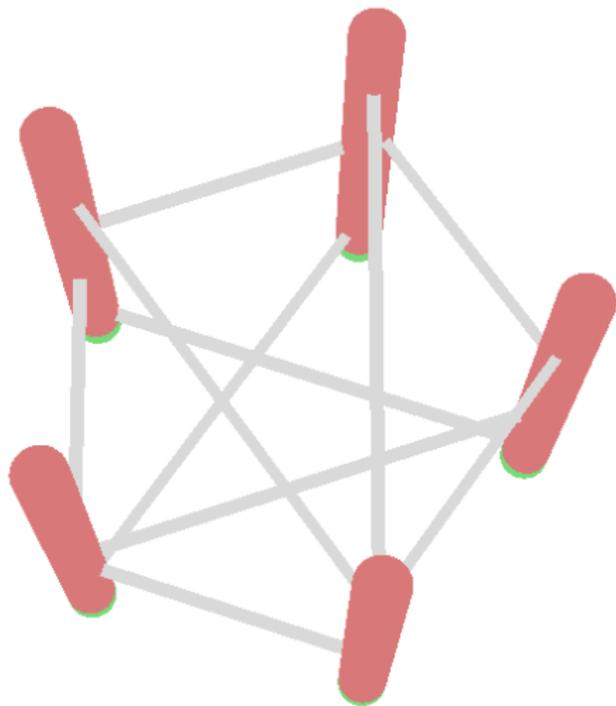
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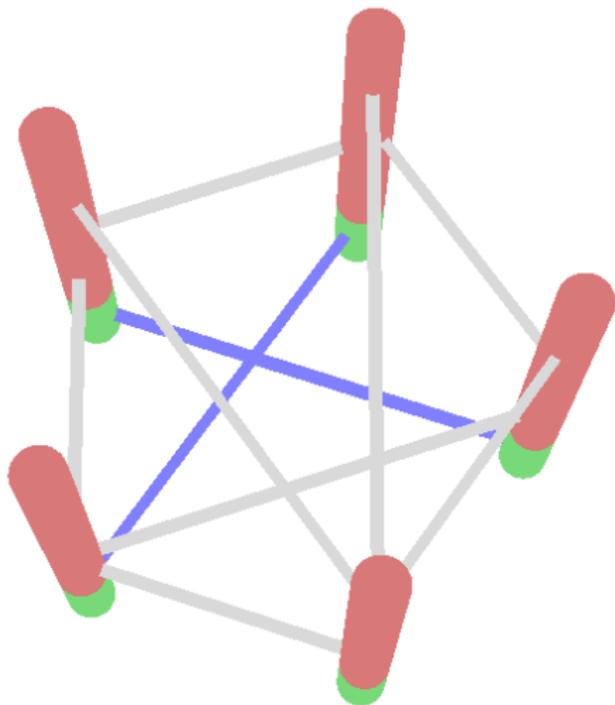
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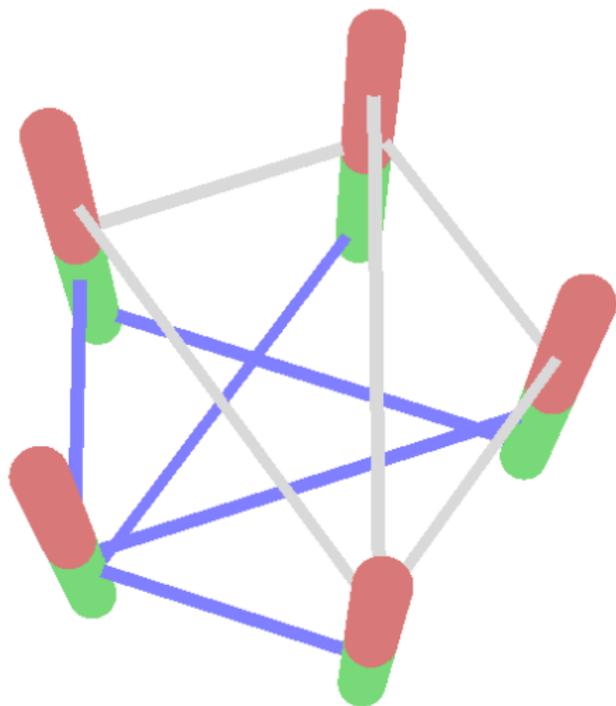
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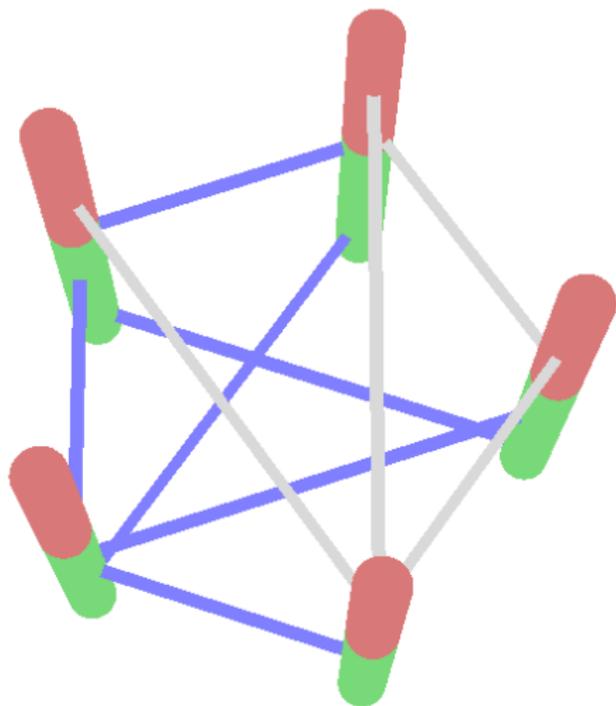
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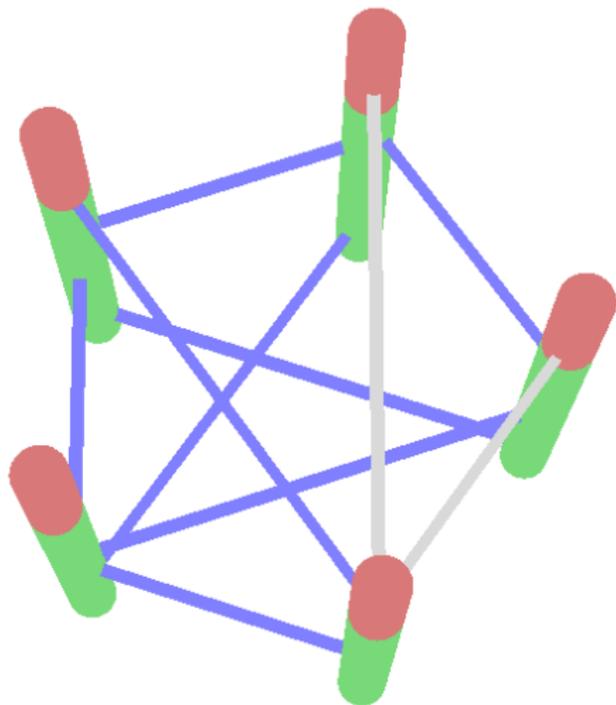
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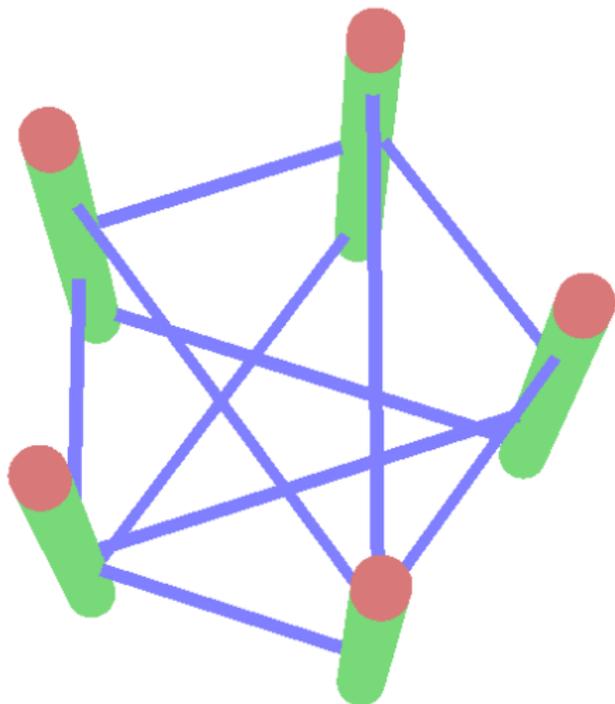
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As in the case of h_1 and h_{conn} both \hat{h}_1 and \hat{h}_{conn} are random variables, but they take values in the interval $[0, 1]$.

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However, the statement that aas $h_1 = h_{\text{conn}}$ is clearly equivalent to the statement that aas $\hat{h}_1 = \hat{h}_{\text{conn}}$.

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In a similar way, for $p_1 \leq p_2$ we have

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THE EVOLUTION OF THE RANDOM GRAPH

If $M = o(\sqrt{n})$ then a.a.s. $G(n, p)$ consists of isolated vertices and isolated edges.

If $M = o(n^{(k-1)/k})$ then a.a.s. all components of $G(n, p)$ are trees with at most k vertices.

If $M = o(n)$ then a.a.s. all components of $G(n, p)$ are trees of size $o(\log n)$.

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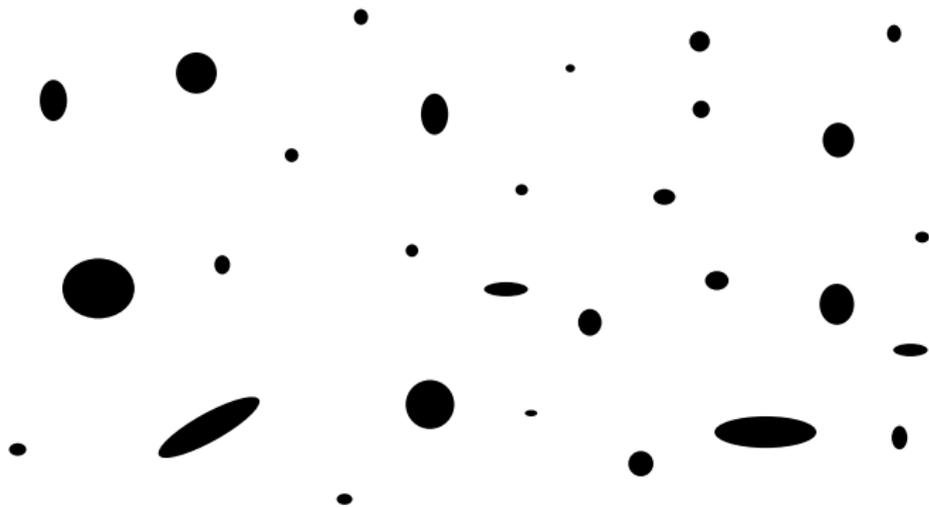
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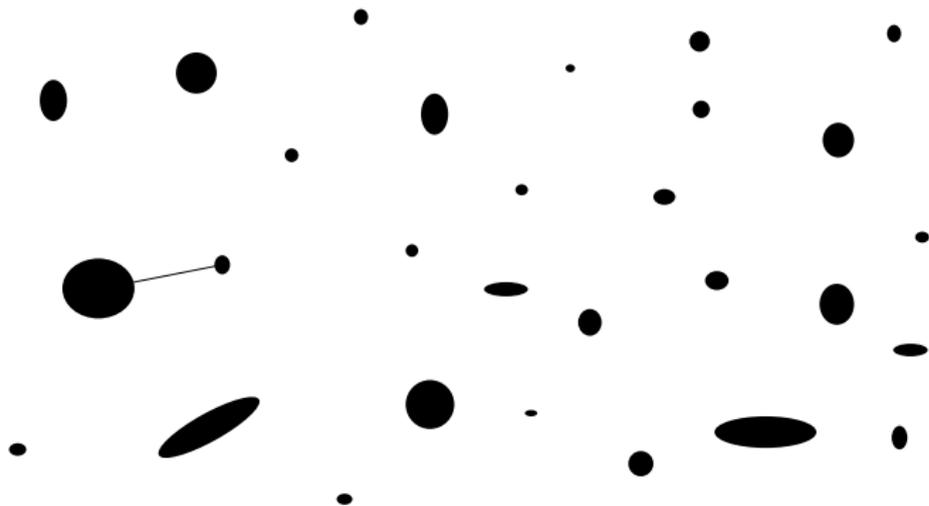
If $M = o(n^{(k-1)/k})$ then a.a.s. all components of $G(n, p)$ are trees with at most k vertices.

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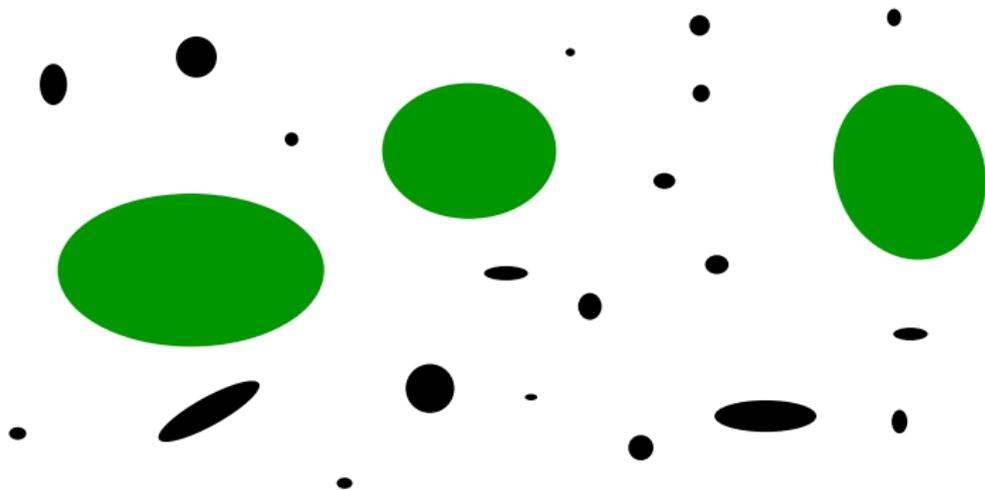
THE SUBCRITICAL PHASE



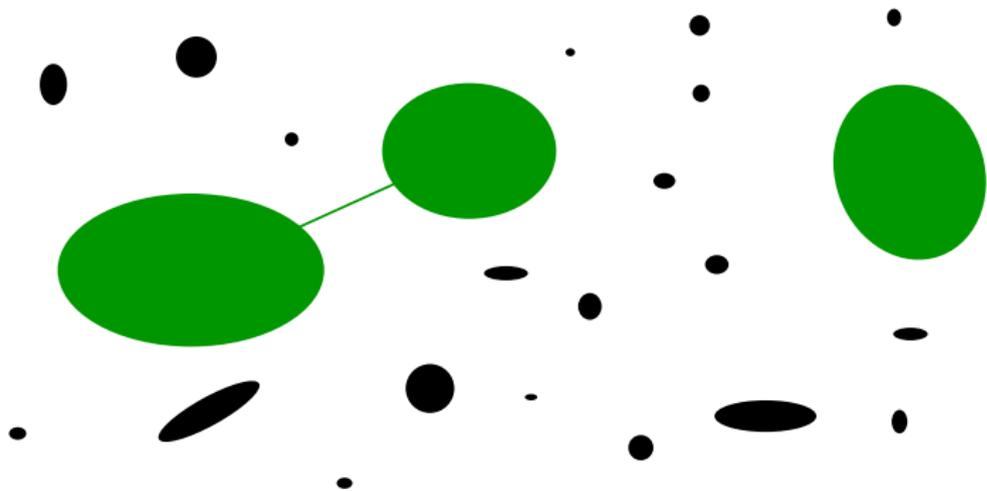
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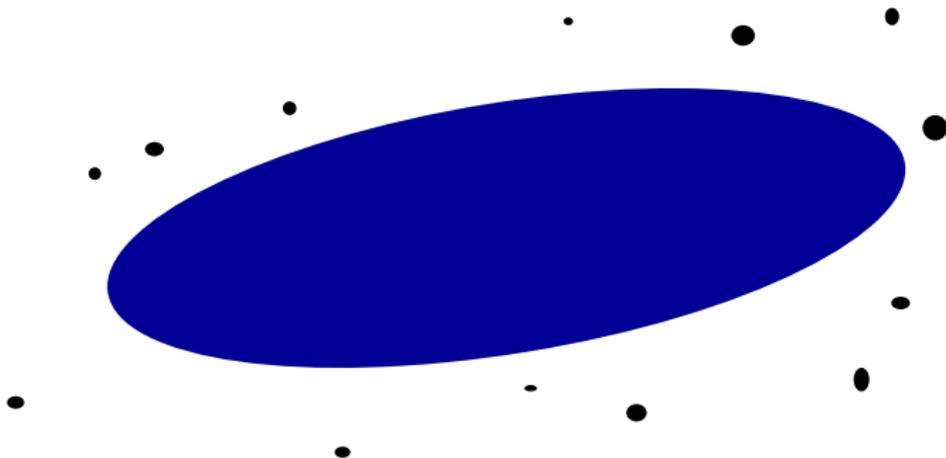
THE CRITICAL PHASE



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THE SUPERCRITICAL PHASE



THE RIGHT SCALING

THEOREM ERDŐS, RÉNYI'60

The “coagulation phase” takes place when $M = (1/2 + o(1))n$.

Thus, for instance, the largest component of $G(n, 0.4999n)$ has $\Theta(\log n)$ vertices, while the size of the largest component of $G(n, 0.5001n)$ is $\Theta(n)$.

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The components start to merge when they are of size $\Theta(n^{2/3})$. It happens when $M = n/2 + \Theta(n^{2/3})$.

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If $np \rightarrow \infty$, then a.a.s $G(n, p)$ contains triangles.

This can be easily proved using the 1st and 2nd moment method we mastered ten minutes ago.

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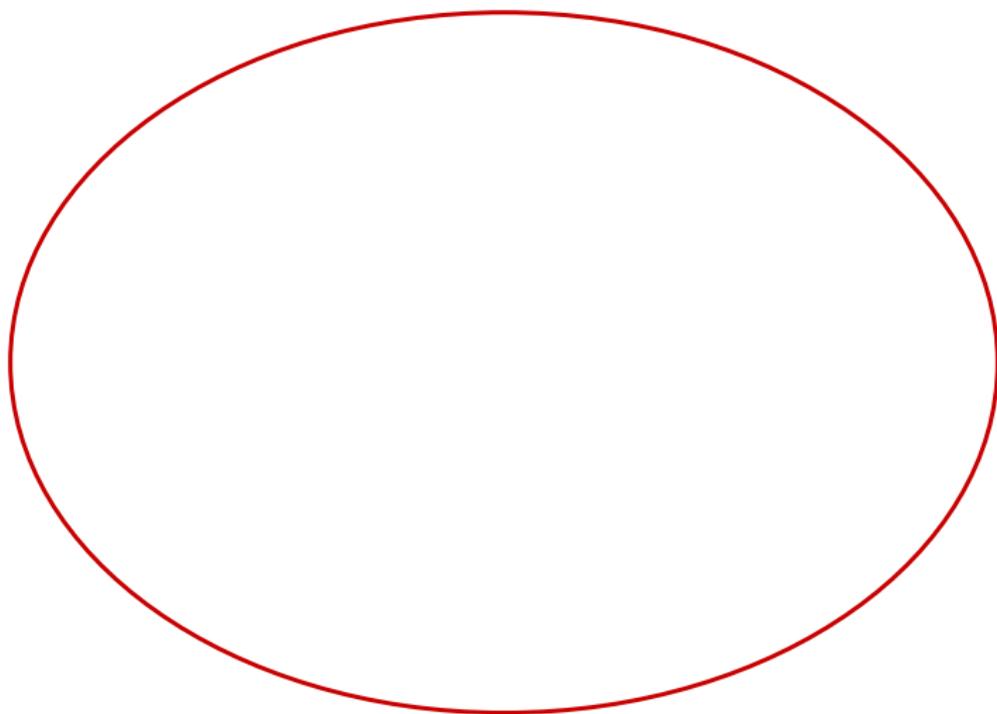
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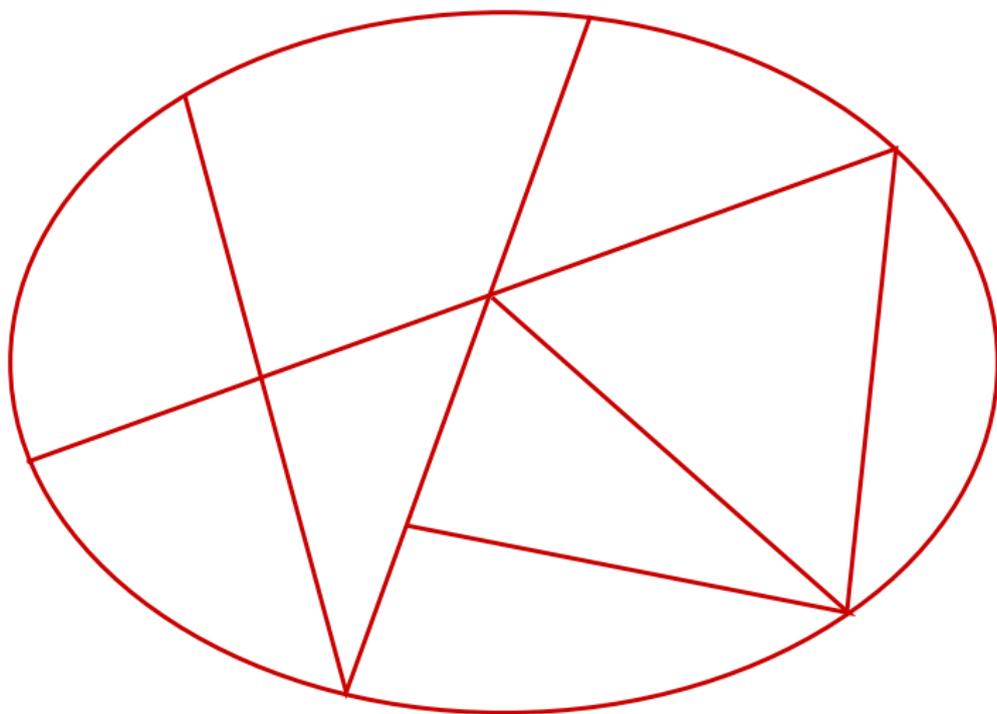
LARGE DEVIATION INEQUALITIES

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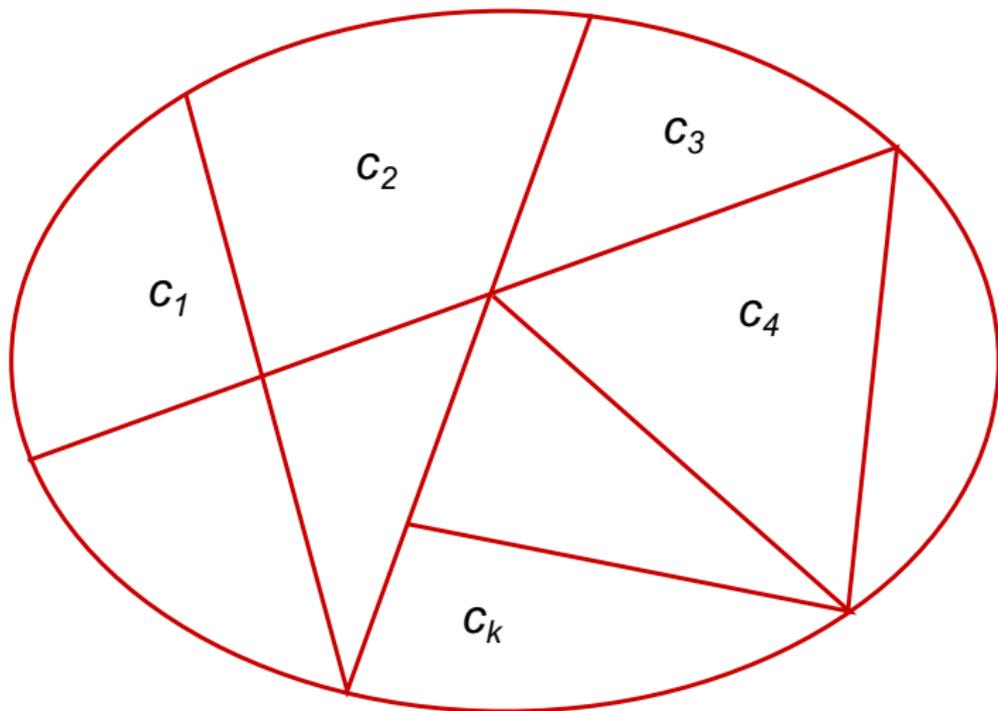
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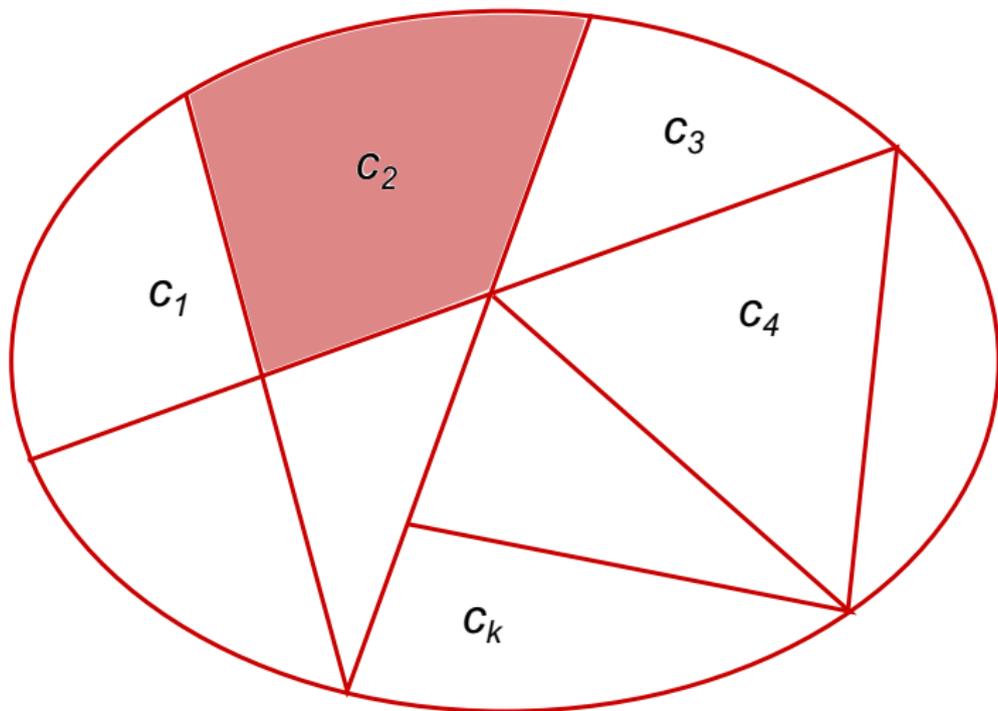
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$$\Pr(|X - \mathbb{E}X| \geq t) \leq 2 \exp\left(-\frac{t^2}{2 \sum_i c_i^2}\right).$$

In particular,

$$\Pr(X = 0) \leq 2 \exp\left(-\frac{(\mathbb{E}X)^2}{2 \sum_i c_i^2}\right).$$

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We want to replace the full sum $\sum_{i=1}^k c_i^2$ by a partial sum of c_i 's.

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TALAGRAND'S INEQUALITY

$$\Pr(|X - \mu X| \geq t) \leq 4 \exp\left(-\frac{t^2}{4w}\right),$$

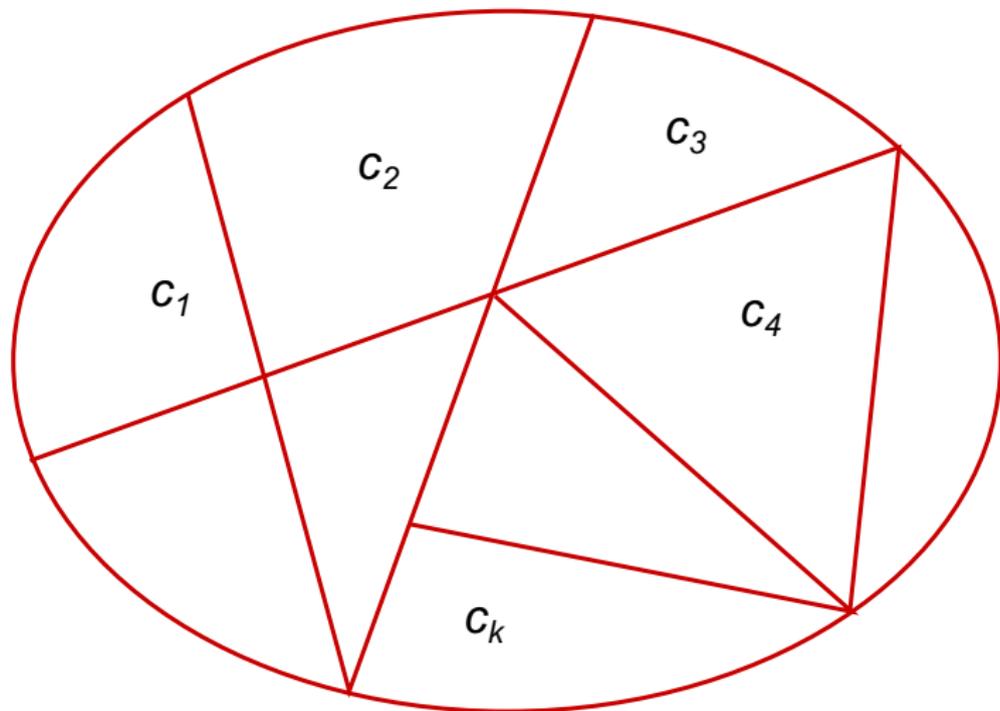
where μX is the median of X and

$$w = \max_{\Lambda} \left\{ \sum_{i \in \Lambda} c_i^2 \right\}$$

where the maximum is taken over all certificates Λ for A .

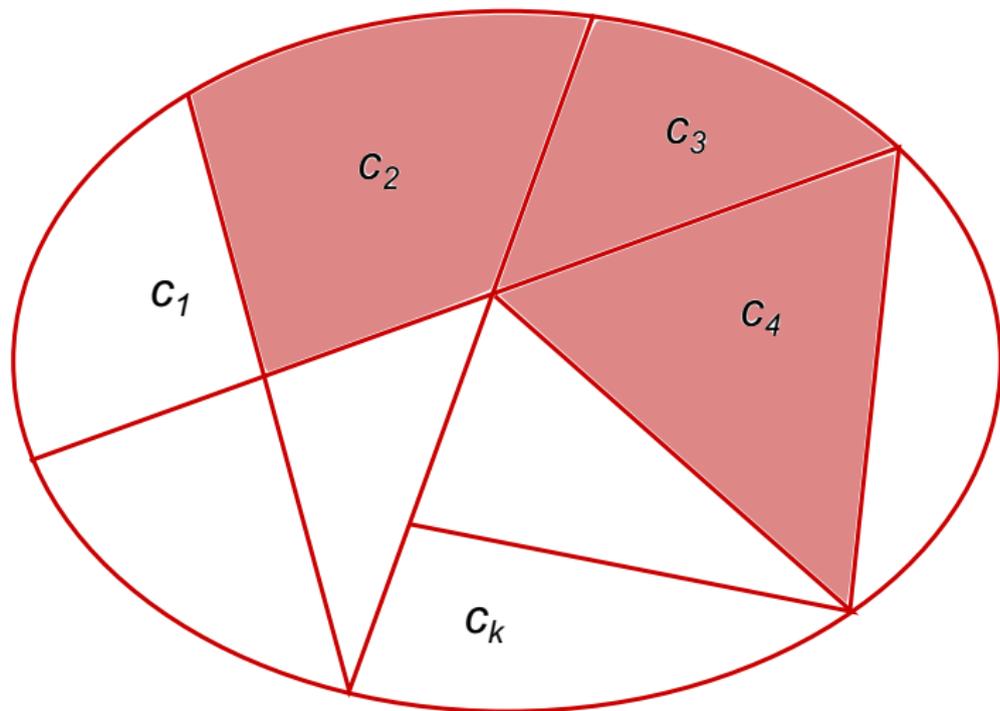
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- (i) In order to certificate that $\alpha(G) \geq r$ it is enough to point out r vertices which belong to this set.
- (ii) There are no small certificates that $\chi(G) \geq r$.
- (iii) The size of the certificate that the number of triangles is larger than r is, of course, $3r$.

THE INDEPENDENCE NUMBER

Let $X = \alpha(G(n, p))$ and $k = 2\mathbb{E}X$. Then random variable $\bar{X} = \min\{X, k\}$ has roughly the same expectation (and median) as X , but its certificate is at most $2\mathbb{E}X$.

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From Azuma's inequality we get

$$\Pr(|X - \mathbb{E}X| \geq t) \leq 2 \exp(-t^2/(2n)),$$

while from Talagrand's inequality, applied to \bar{X} , we get roughly

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In particular, for every $\gamma \rightarrow \infty$,

$$\Pr(|X - \mathbb{E}X| \geq \gamma\sqrt{\mathbb{E}X}) \rightarrow 0.$$

THE PROBABILITY THAT THERE ARE NO TRIANGLES

Let X denote the number of triangles in $G(n, p)$ and \bar{X} be the maximum size of the family of edge-disjoint triangles.

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From Talagrand's inequality we get

$$\begin{aligned}\Pr(X = 0) &= \Pr(\hat{X} = 0) \leq \Pr(|\hat{X} - \mathbb{E}\hat{X}| \geq \mathbb{E}\hat{X}) \\ &\leq 4 \exp\left(-\frac{(\mathbb{E}\hat{X})^2}{12\mathbb{E}X}\right) \leq 4 \exp\left(-\frac{\mathbb{E}X}{108}\right).\end{aligned}$$

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On the other hand, from FKG inequality we get

$$\begin{aligned} \Pr(X = 0) &\geq (1 - p^3)^{\binom{n}{3}} = e^{-(1+o(1))\binom{n}{3}p^3} \\ &= \exp(-(1 + o(1))\mathbb{E}X). \end{aligned}$$

REMARKS

THEOREM JANSON, ŁUCZAK, RUCIŃSKI '90

Let $X(H)$ count the number of copies of H in $G(n, p)$. Then, for every H , we have

$$\Pr(X(H) = 0) = \exp\left(-\Theta\left(\min_{F \subseteq H} \mathbb{E}X(F)\right)\right).$$

Although we know that

$$\Pr(X(K_3) = 0) = \exp\left(-\Theta\left(\min\{\mathbb{E}X(K_3), \mathbb{E}X(K_2)\}\right)\right),$$

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The first factor can be bounded from above by $\exp(-cM)$, the second one, by our theorem and the equivalence results, is smaller than $\exp(-c'M)$ and it turns out that $c' > c$. Hence $\mathbb{E}Y \rightarrow 0$ and the assertion follows from the first moment method. □

MAKER-BREAKER GAME $MB(n, q, H)$

Two players: **Maker** and **Breaker**

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Maker wins if **his graph** contains a copy of H
otherwise the win comes to Breaker.

THRESHOLD BIAS

The **threshold bias** $\bar{q}(n) = \bar{q}_{\mathcal{A}}(n)$ is the maximum q so that Maker can win $MB(n, q, \mathcal{A})$.

i.e. Maker has a winning strategy to build a graph with $\binom{n}{2} / (q + 1)$ edges which has property \mathcal{A} .

$MB(n, q, K_3)$

CLAIM FOLKLORE

In $MB(n, q, K_3)$, when Maker tries to build a triangle, the threshold bias is $\Theta(\sqrt{n})$.

More specifically:

- ▶ Maker has a winning strategy if $q < \sqrt{n}$,
- ▶ Breaker has a winning strategy if $q > 2\sqrt{n}$.

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The threshold bias for $MB(n, q, K_3)$ lies in the interval $[\sqrt{n}, 2\sqrt{n}]$.

We aim into the following exciting result.

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I can understand it...

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but you should know that the method we shall present (introduced by **BEDNARSKA, ŁUCZAK'99**) is the only known method which gives the right order of bias for every H !

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The edges chosen by Maker form a graph $\hat{F} = G(n, M)$, with $M = n^{3/2}$.

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However, not every such an edge is in his graph – because of his strategy, some of the edges he selects has already been claimed by Breaker and so they are ‘lost’ and will not belong to \hat{F} .

However, since the choice is random, with a very large probability fewer than 1% of edges of $\hat{F} = G(n, M)$ have been claimed by Breaker, i.e. more than 99% of edges of \hat{F} are in Maker’s graph!

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But is this the end of the proof?

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But we know that $G(n, M)$ has the property that it contains a triangle in every subgraph which have at least $0.99M$ edges!
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PROOF

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This is the end (ADELE'12)!

Since only one of the player can have a winning strategy, if Maker has got a strategy that wins sometimes, he has also got a strategy which wins always (since Breaker cannot have it). □

THE INDEPENDENCE NUMBER

PROBLEM

What is the independence number of $G(n, p)$, say, for $p = \log n/n$?

FACT

Let $p = \log n/n$, $\epsilon > 0$ and $k = n \log \log n / \log n$. Then, as $n \rightarrow \infty$,
 $\alpha(G(n, p)) \leq (2 + \epsilon)k$.

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Proof The first moment method. Estimate $\mathbb{E}X$, where X is the number of independent subsets of size $(2 + \epsilon)k$. Then

$$\mathbb{E}X = \binom{n}{(2 + \epsilon)k} (1 - p)^{\binom{(2 + \epsilon)k}{2}} \rightarrow 0.$$

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Proof Surprisingly, this result can also be proved by the first moment method. Estimate $\mathbb{E}Y$, where Y is the number of covering subsets of size $(1 - \epsilon)k$. Then

$$\mathbb{E}Y = \binom{n}{(1 - \epsilon)k} (1 - (1 - p)^{(1 - \epsilon)k})^{n - (1 - \epsilon)k} \rightarrow 0.$$

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Let $p = \log n/n$, $\epsilon > 0$ and $k = n \log \log n / \log n$. Then, aas

$$(1 - \epsilon)k \leq \alpha(G(n, p)) \leq (2 + \epsilon)k .$$

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$$P(|\alpha(G(n, p)) - \mathbb{E}\alpha(G(n, p))| \geq t) \leq 4 \exp(-t^2/9k) .$$

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Let X count independent sets of size $(2 - \epsilon)k$.

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Two random sets of this size share $\Theta(k^2/n)$ vertices, so we cannot expect that the existence of one set in such a pair is “almost independent” from the existence of the second one. After some (fairly long) calculations one can show that

$$\mathbb{E}X(X - 1) \geq (\mathbb{E}X)^2 \exp\left(\frac{2k}{(\log \log n)^3}\right) .$$

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$$\Pr(X = 0) \leq \frac{\text{Var}X}{(\mathbb{E}X)^2} \quad \text{but} \quad \frac{\text{Var}X}{(\mathbb{E}X)^2} \gg 1 \quad (\text{sic!})$$

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$$\Pr(X > 0) \geq \frac{(\mathbb{E}X)^2}{\mathbb{E}X^2} \geq \exp\left(-\frac{3k}{(\log \log n)^3}\right)$$

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It seems that the 2nd moment method is completely useless in this case!

FRIEZE'S IDEA: COMBINE CAUCHY AND TALAGRAND!

The main idea of Frieze's argument

We want to show that as $n \rightarrow \infty$, $\alpha(G(n, p)) \geq (2 - 3\epsilon)k$.

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Let us assume that this is not the case, i.e. that

$$\mathbb{E}\alpha(G(n, p)) \leq (2 - 2\epsilon)k$$

and hope to get a contradiction.

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This is the contradiction we have been hoping for!



TRIANGLES: SOME FURTHER REMARKS

(EASY) COROLLARY OF LARGE DEVIATION INEQUALITIES

If $M = n^{3/2}$, then as we cannot destroy all triangles in $G(n, M)$ by removing $0.01M$ edges.

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Here is a much harder result.

THEOREM HAXELL, KOHAYAKAWA, ŁUCZAK'96

If $M = n^{3/2}$, then as we cannot destroy all triangles in $G(n, M)$ by removing $0.49M$ edges.

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or one of the transferring theorems (by CONLON, GOWERS and SCHACHT)
or hypergraph containers (by SAXTON, THOMASSON and BALOGH, MORRIS, SAMOTIJ).

ALTHOUGH THIS TALK WAS BROUGHT TO YOU COMPLETELY COMMERCIAL-FREE...

THEOREM ERDŐS, RÉNYI'60

If $np \rightarrow 0$, then aas $G(n, p)$ contains no triangles.

If $np \rightarrow \infty$, then aas $G(n, p)$ contains triangles.

THEOREM ERDŐS, RÉNYI'59

Let $p(n) = \frac{1}{n}(\ln n + \gamma(n))$. Then

$$\lim_{n \rightarrow \infty} \Pr(G(n, p) \text{ is connected}) = \begin{cases} 0 & \text{if } \gamma(n) \rightarrow -\infty, \\ 1 & \text{if } \gamma(n) \rightarrow \infty. \end{cases}$$

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We say that the property " $G \not\supseteq K_3$ " has a coarse threshold, while the property " G is connected" has a sharp threshold.

THRESHOLDS

PROBLEM

Can we (combinatorially) characterize graph properties which have sharp thresholds?

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but there exists a nice application to random groups.

Thank you!

FURTHER READINGS

If you are interested in the subject, there are three books on random graphs you might want to read.

B. Bollobás, *Random graphs*, Cambridge University Press, 2nd edition, 2011.

S. Janson, T. Łuczak, A. Ruciński, *Random graphs*, Wiley, 2000.

A. Frieze, M. Karoński, *Introduction to random graphs*, Cambridge University Press, to be published this year.