

# New sendfile(2)

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# Miserable life w/o sendfile(2)

```
while ((cnt = read(filefd, buf, (u_int)blksize))  
    write(netfd, buf, cnt) == cnt)  
    byte_count += cnt;
```

send\_data() в src/libexec/ftpd/ftpd.c,  
FreeBSD 1.0, 1993



# sendfile(2) introduced

```
int  
sendfile(int fd, int s, off_t offset, size_t nbytes, .. );
```

- 1997: HP-UX 11.00
- 1998: FreeBSD 3.0 and Linux 2.2



# sendfile(2) in FreeBSD

- First implementation - mapping userland cycle to the kernel:
  - `read(filefd)` → `VOP_READ(vnode)`
  - `write(netfd)` → `sosend(socket)`
  - `blksize` → `PAGE_SIZE`



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- Further optimisations:
  - 2004: `SF_NODISKIO` flag
  - 2006: inner cycle, working on `sbspace()` bytes
  - 2013: sending a shared memory descriptor data



# Problem #1: blocking on I/O

Algorithm of a modern HTTP-server:

- ① Take yet another descriptor from kevent(2)
- ② Do write(2)/read(2)/sendfile(2) on it
- ③ Go to 1



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Algorithm of a modern HTTP-server:

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Bottleneck: any syscall time.



# Attempts to solve problem #1

- Separate I/O contexts: processes, threads
  - Apache
  - nginx 2



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- Separate I/O contexts: processes, threads
  - Apache
  - nginx 2
- SF\_NODISKIO + aio\_read(2)
  - nginx
  - Varnish



# More attempts . . .

- aio\_mlock(2) instead of aio\_read(2)
- aio\_sendfile(2) ???



## Problem #2: control over VM

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- VOP\_READ() leaves pages in VM cache
- VOP\_READ() [for UFS] does readahead
- Not easy to prevent it doing that!



# waht if VOP\_GETPAGES()?

~~VOP\_READ()~~ → VOP\_GETPAGES()

- Pros:

- sendfile() already works on pages
- implementations for vnode and shmem converge
- control over VM is now easier task



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- Cons
  - Losing readahead heuristics ☹



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- Pros:
  - sendfile() already works on pages
  - implementations for vnode and shmem converge
  - control over VM is now easier task
- Cons
  - Losing readahead heuristics ☹
  - But no one used them! ☺



# VOP\_GETPAGES\_ASYNC()

*int*

```
VOP_GETPAGES(struct vnode *vp, vm_page_t *ma,  
int count, int reqpage);
```

- ① Initialize buf(9)
- ② buf->b\_iodone = bdone;
- ③ bstrategy(buf);
- ④ bwait(buf); /\* sleeps until I/O completes \*/
- ⑤ return;



# VOP\_GETPAGES\_ASYNC()

```
int  
VOP_GETPAGES_ASYNC(struct vnode *vp,  
vm_page_t *ma, int count, int reqpage,  
vop_getpages_iodone_t *iodone, void *arg);
```

- ① Initialize buf(9)
- ② buf->b\_iodone = vnode\_pager\_async\_iodone;
- ③ bstrategy(buf);
- ④ return;

vnode\_pager\_async\_iodone calls **iodone()** .



# naive non-blocking sendfile(2)

In kern\_sendfile():

- ① nios++;
- ② VOP\_GETPAGES\_ASYNC(sendfile\_iodone);

In sendfile\_iodone():

- ① nios--;
- ② if (nios) return;
- ③ sosend();



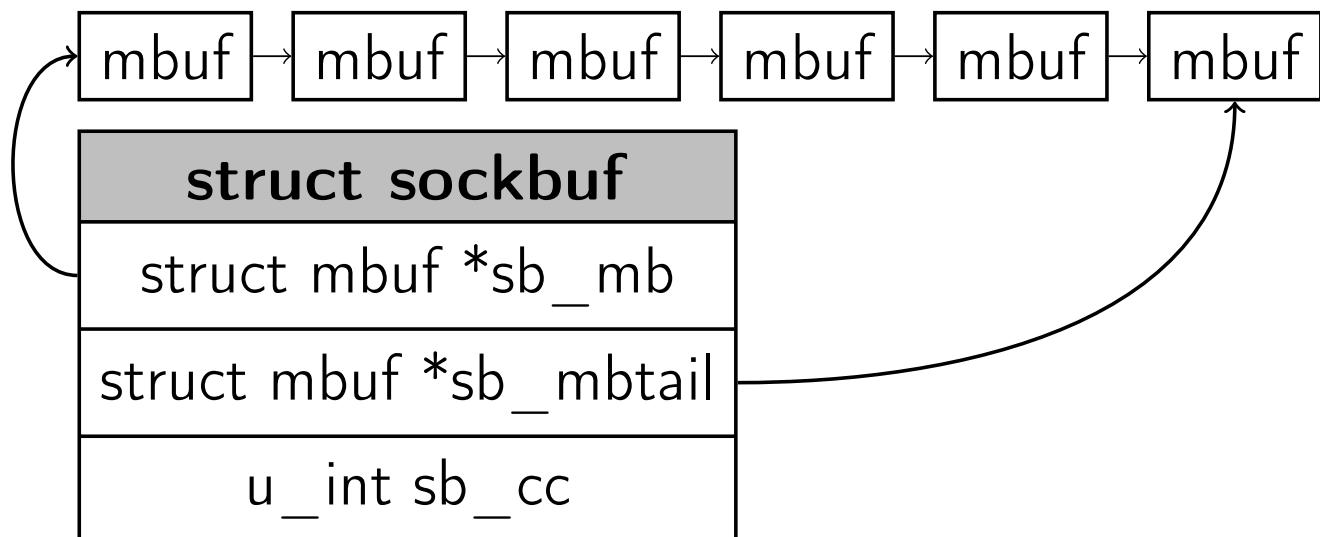
# the problem of naive implementation

```
sendfile(filefd, sockfd, ...);  
write(sockfd, ...);
```

New sendfile(2) "not ready" data in socket buffers

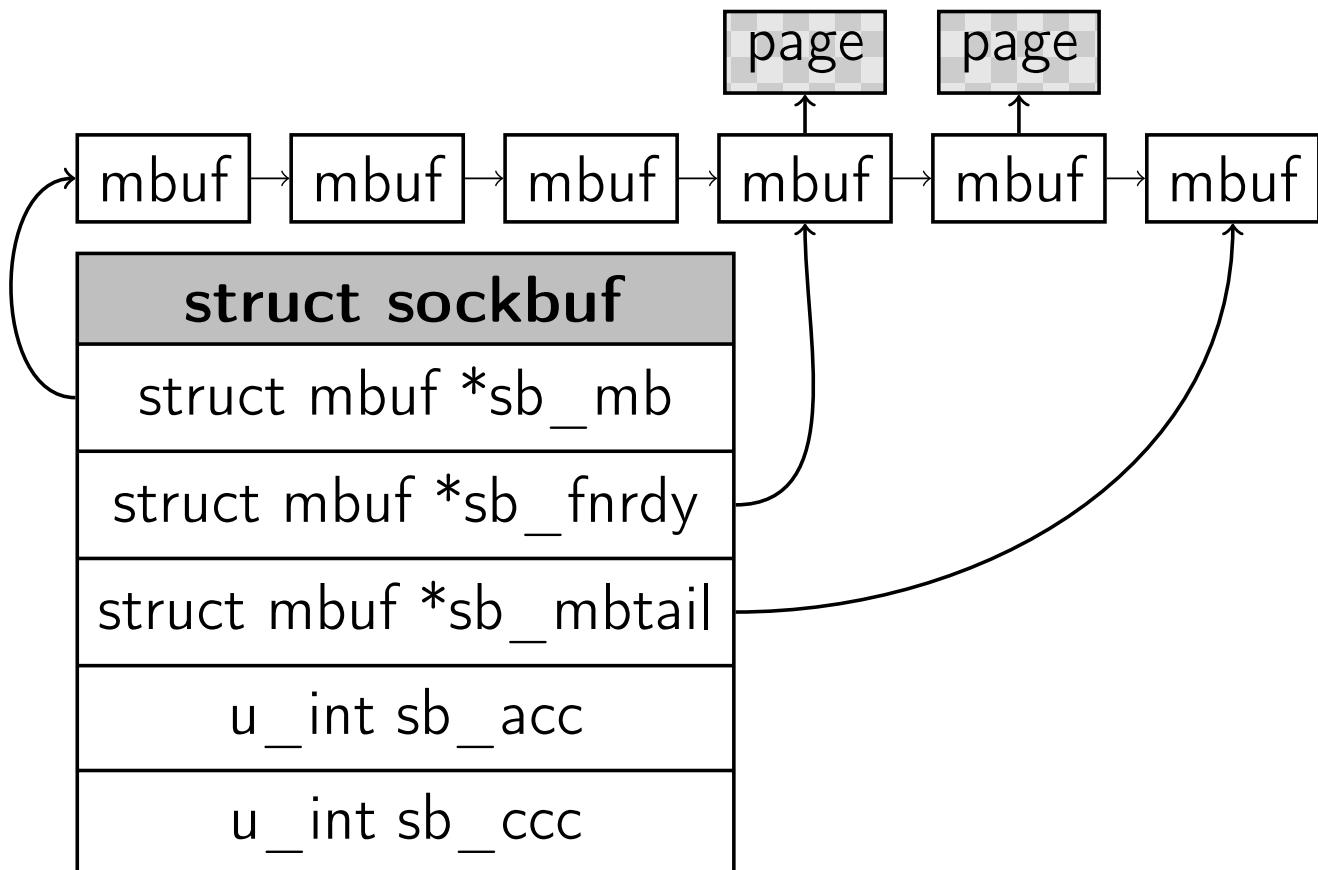


# socket buffer





# socket buffer with “not ready” data





# non-blocking sendfile(2)

In kern\_sendfile():

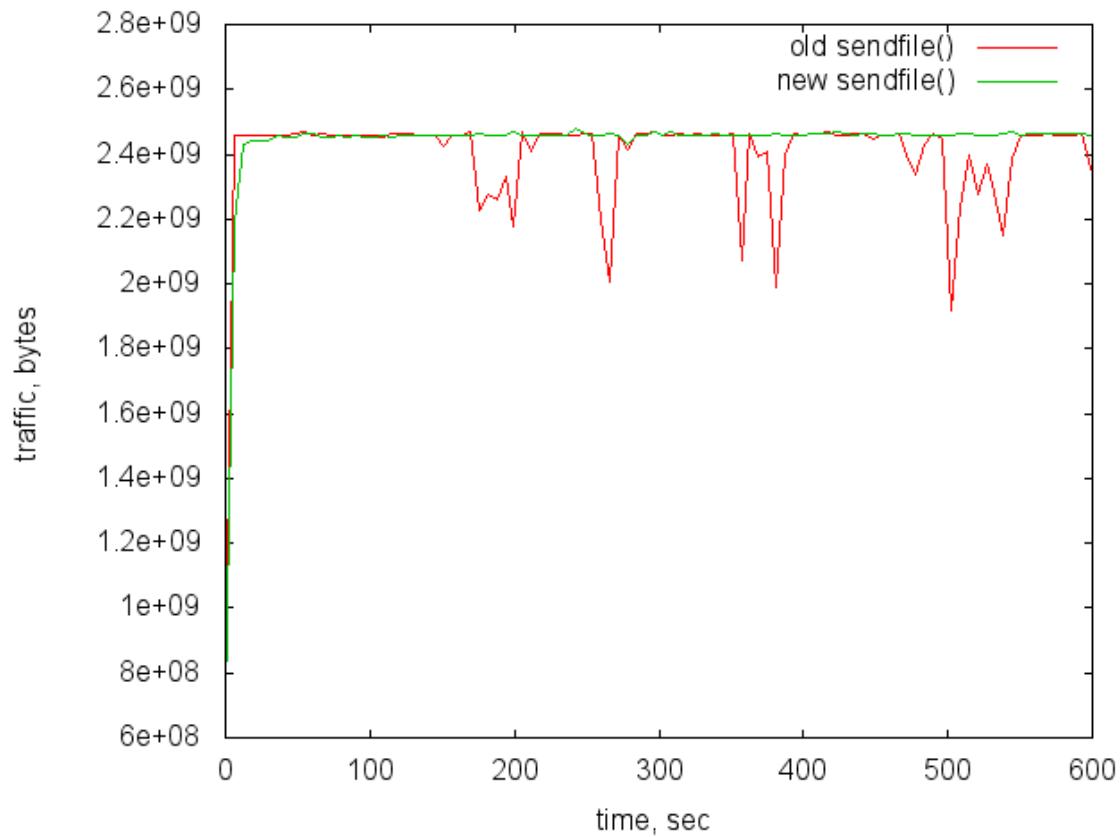
- ① nios++;
- ② VOP\_GETPAGES\_ASYNC(sendfile\_iodone);
- ③ sosend(NOT\_READY);

In sendfile\_iodone():

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- ② if (nios) return;
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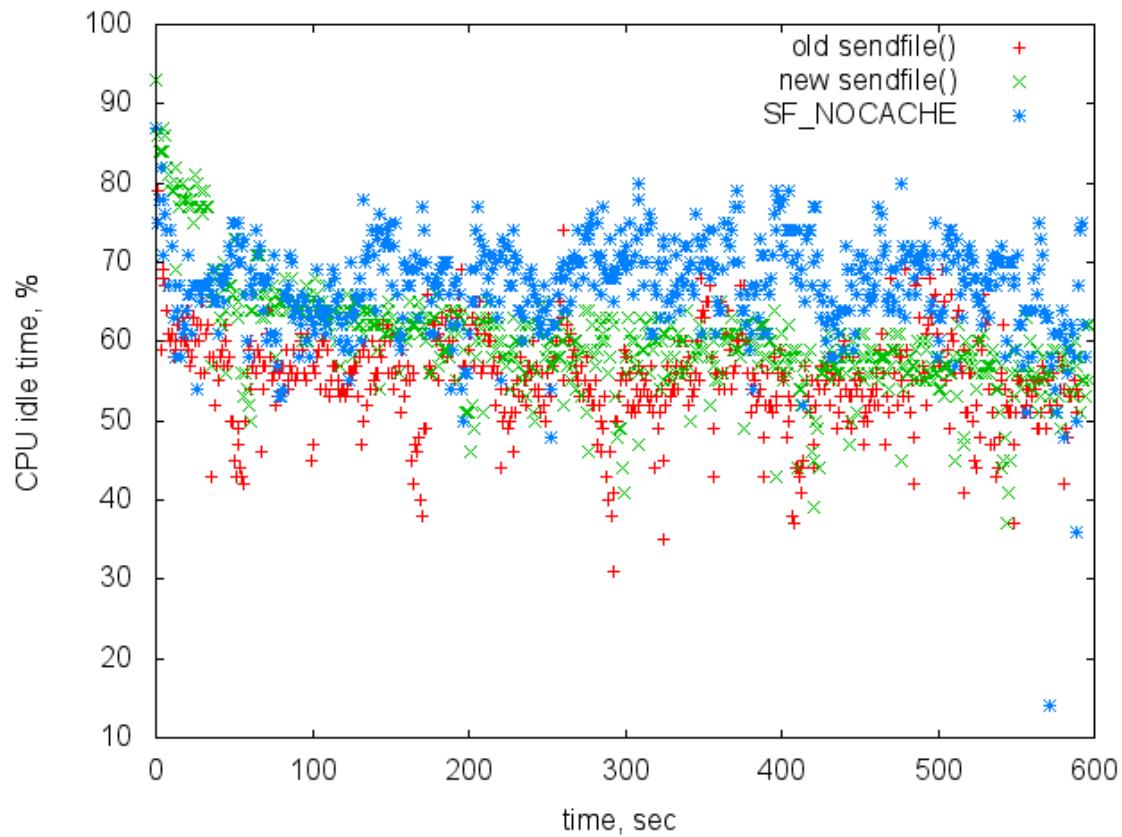


## traffic





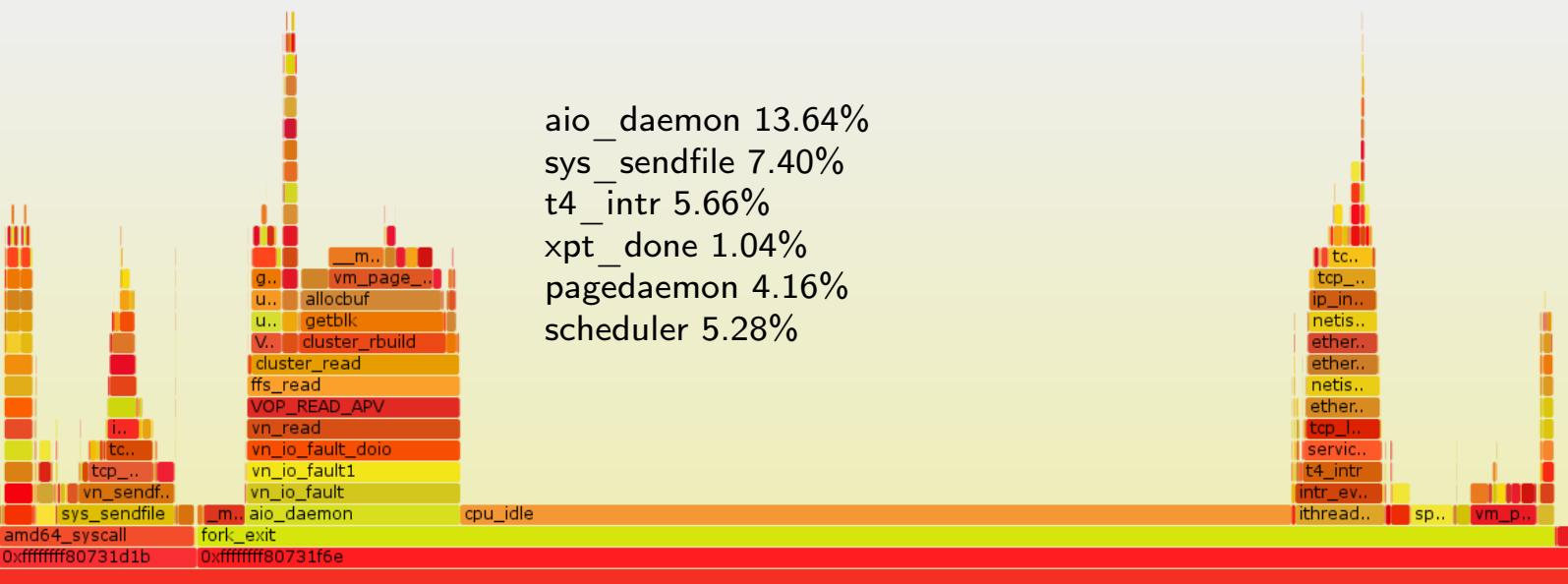
# CPU idle





# profiling sendfile(2) in head

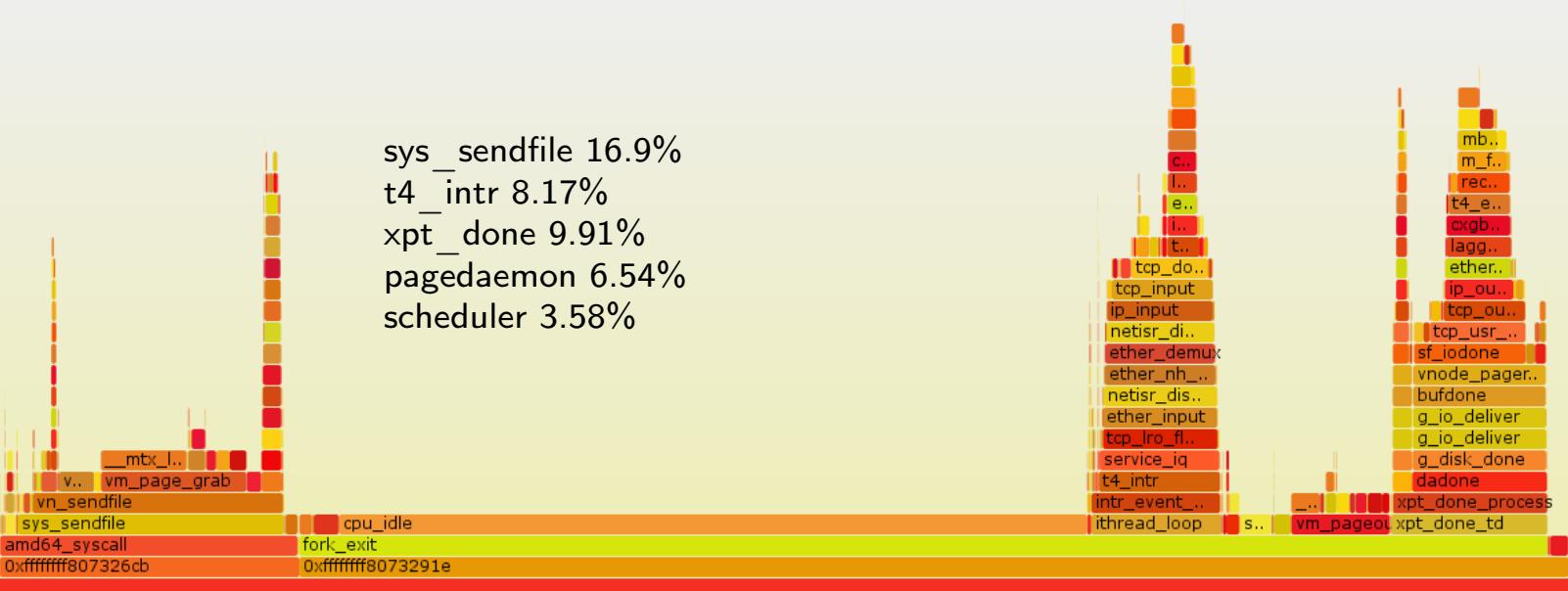
aio\_daemon 13.64%  
 sys\_sendfile 7.40%  
 t4\_intr 5.66%  
 xpt\_done 1.04%  
 pagedaemon 4.16%  
 scheduler 5.28%





# profiling new sendfile(2)

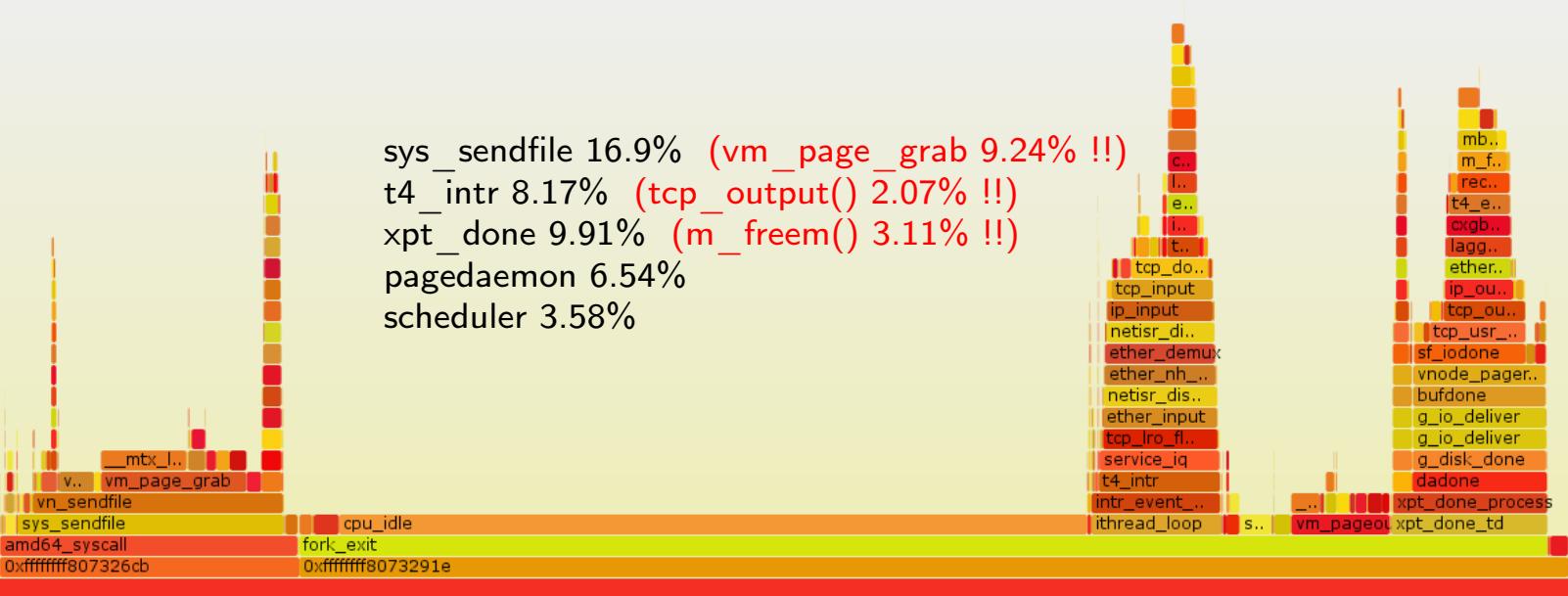
sys\_sendfile 16.9%  
 t4\_intr 8.17%  
 xpt\_done 9.91%  
 pagedaemon 6.54%  
 scheduler 3.58%





# profiling new sendfile(2)

sys\_sendfile 16.9% (vm\_page\_grab 9.24% !!)  
 t4\_intr 8.17% (tcp\_output() 2.07% !!)  
 xpt\_done 9.91% (m\_freem() 3.11% !!)  
 pagedaemon 6.54%  
 scheduler 3.58%





# what did change?

- New code always sends full socket buffer
  - Which is good for TCP (as protocol)
  - Which hurts VM, mbuf allocator,  
and unexpectedly TCP stack

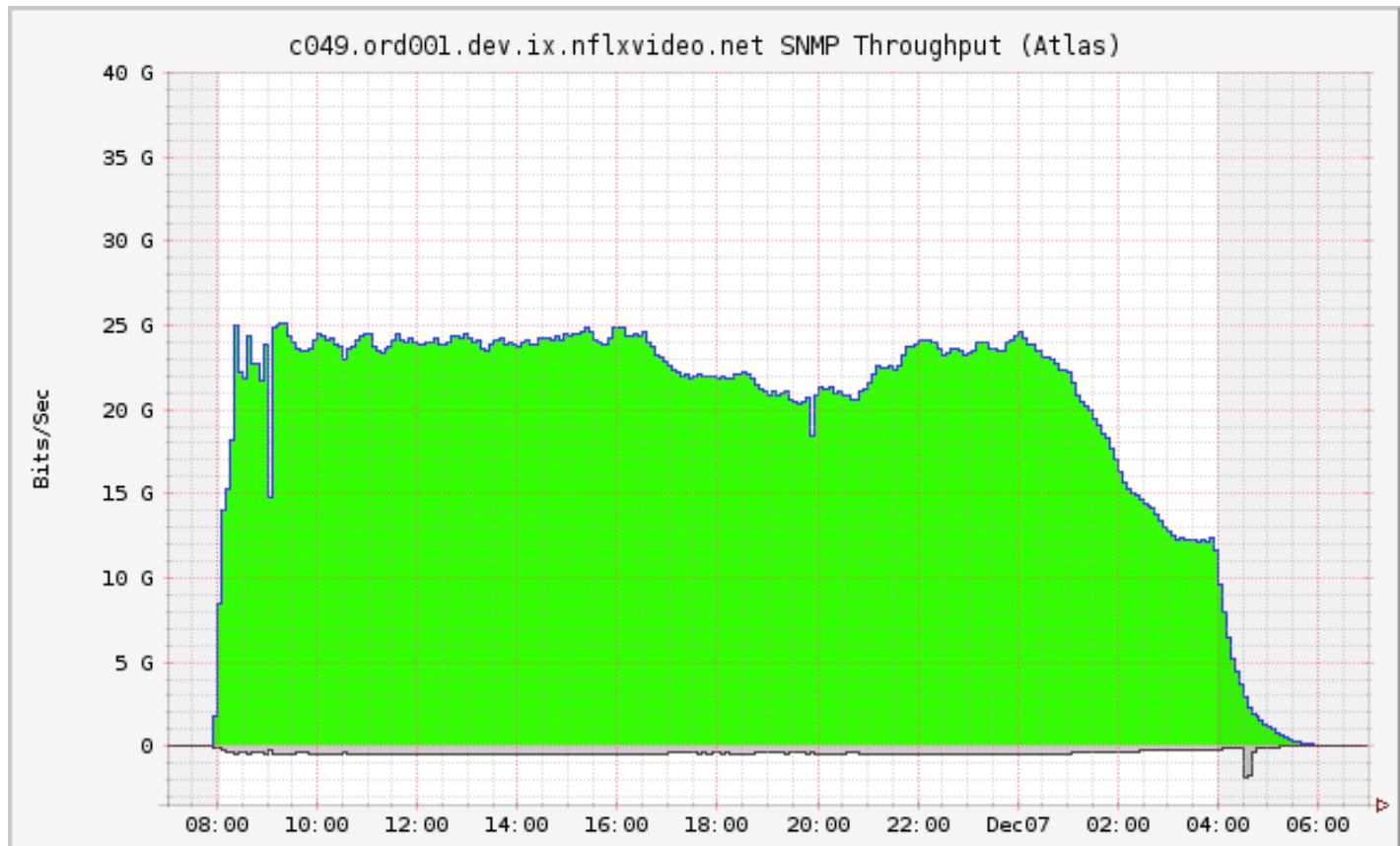


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- Will fix that!

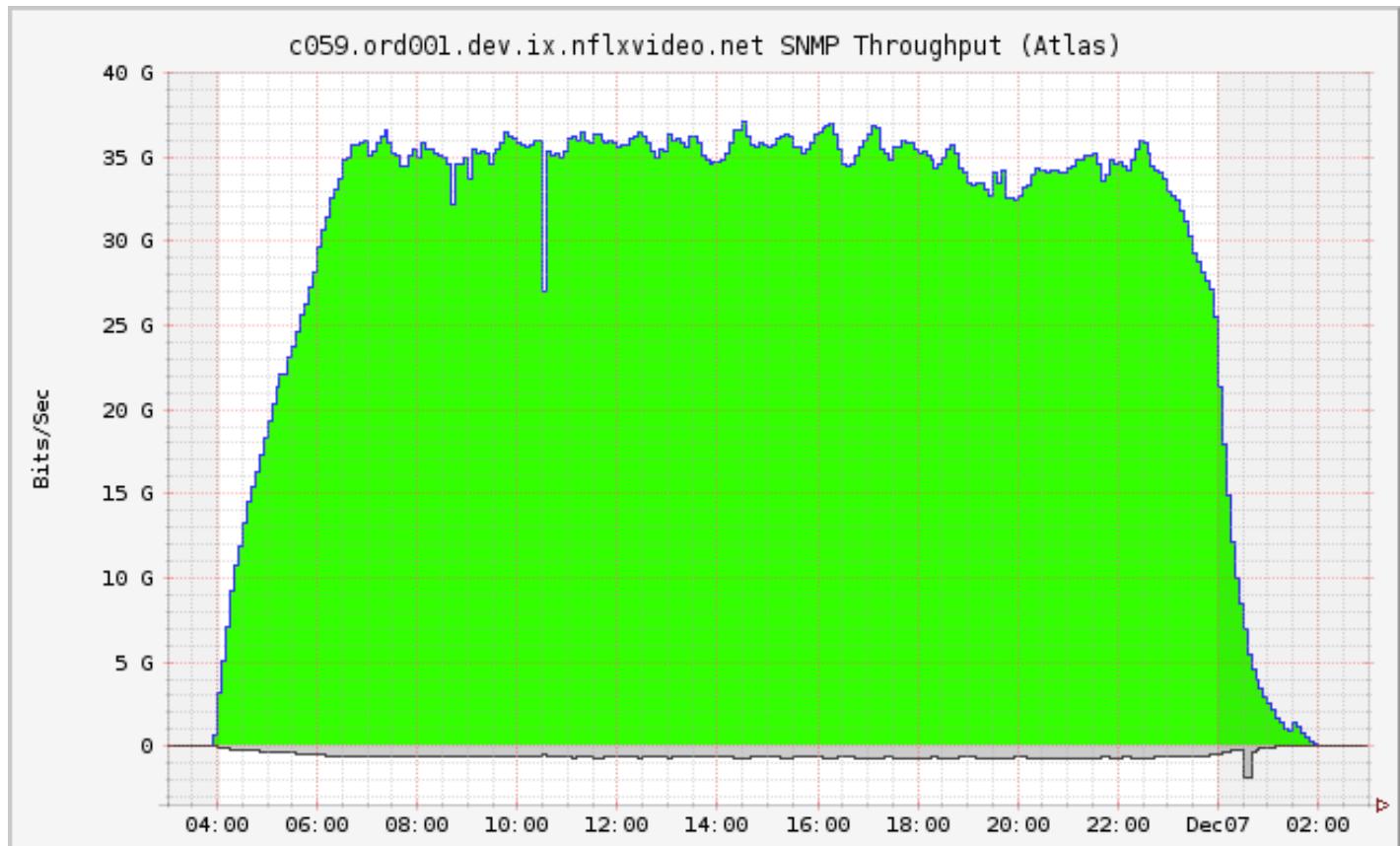


# old sendfile(2) @ Netflix





# new sendfile(2) @ Netflix





# TODO list

## Problems:

- VM & I/O overcommit
- ZFS
- SCTP



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## Future plans:

- sendfile(2) doing TLS

# Questions?