Windows Kernel Internals
Overview

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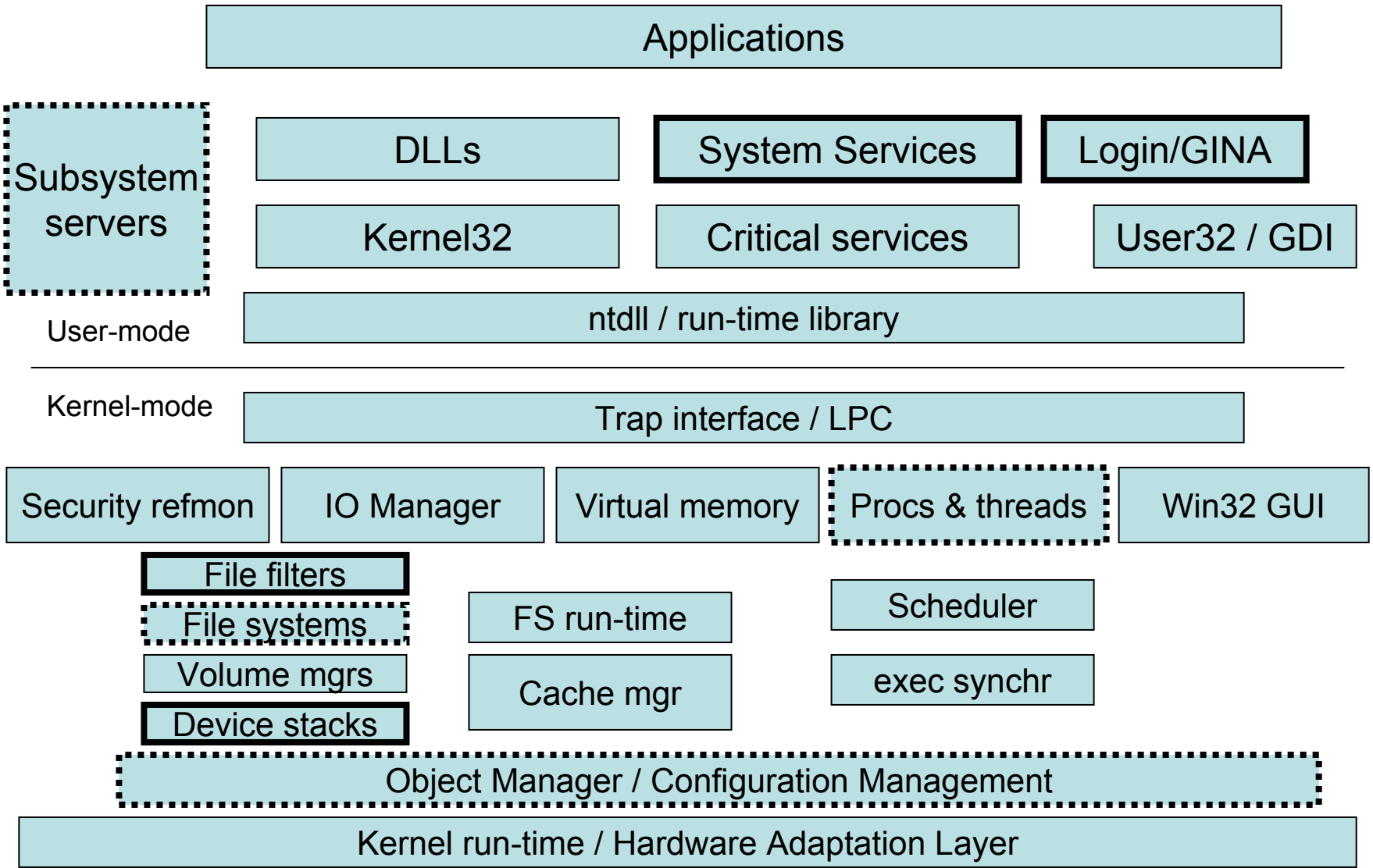
Windows History

- Team formed in November 1988
- Less than 20 people
- Build from the ground up
  - Advanced Operating System
  - Designed for desktops and servers
  - Secure, scalable SMP design
  - All new code
- Rigorous discipline – developers wrote very detailed design docs, reviewed/discussed each others docs and wrote unit tests

Goals of the NT System

- Reliability – Nothing should be able to crash the OS. Anything that crashes the OS is a bug and we won't ship until it is fixed
- Security – Built into the design from day one
- Portability – Support more than one processor, avoid assembler, abstract HW dependencies.
- Extensibility – Ability to extend the OS over time
- Compatibility – Apps must run
- Performance – All of the above are more important than raw speed!

Windows Architecture



Windows Kernel Organization

Kernel-mode organized into

- NTOS (kernel-mode services)
  - Run-time Library, Scheduling, Executive services, object manager, services for I/O, memory, processes, …
- Hal (hardware-adaptation layer)
  - Insulates NTOS & drivers from hardware dependencies
  - Providers facilities, such as device access, timers, interrupt servicing, clocks, spinlocks
- Drivers
  - kernel extensions (primarily for device access)

Major Kernel Services

Process management

Process/thread creation

Security reference monitor

Access checks, token management

Memory manager

Pagefaults, virtual address, physical frame, and pagefile management

Services for sharing, copy-on-write, mapped files, GC support, large apps

Lightweight Procedure Call (LPC)

Native transport for RPC and user-mode system services.

I/O manager (& plug-and-play & power)

Maps user requests into IRP requests, configures/manages I/O devices, implements services for drivers

Cache manager

Provides file-based caching for buffer file system I/O

Built over the memory manager

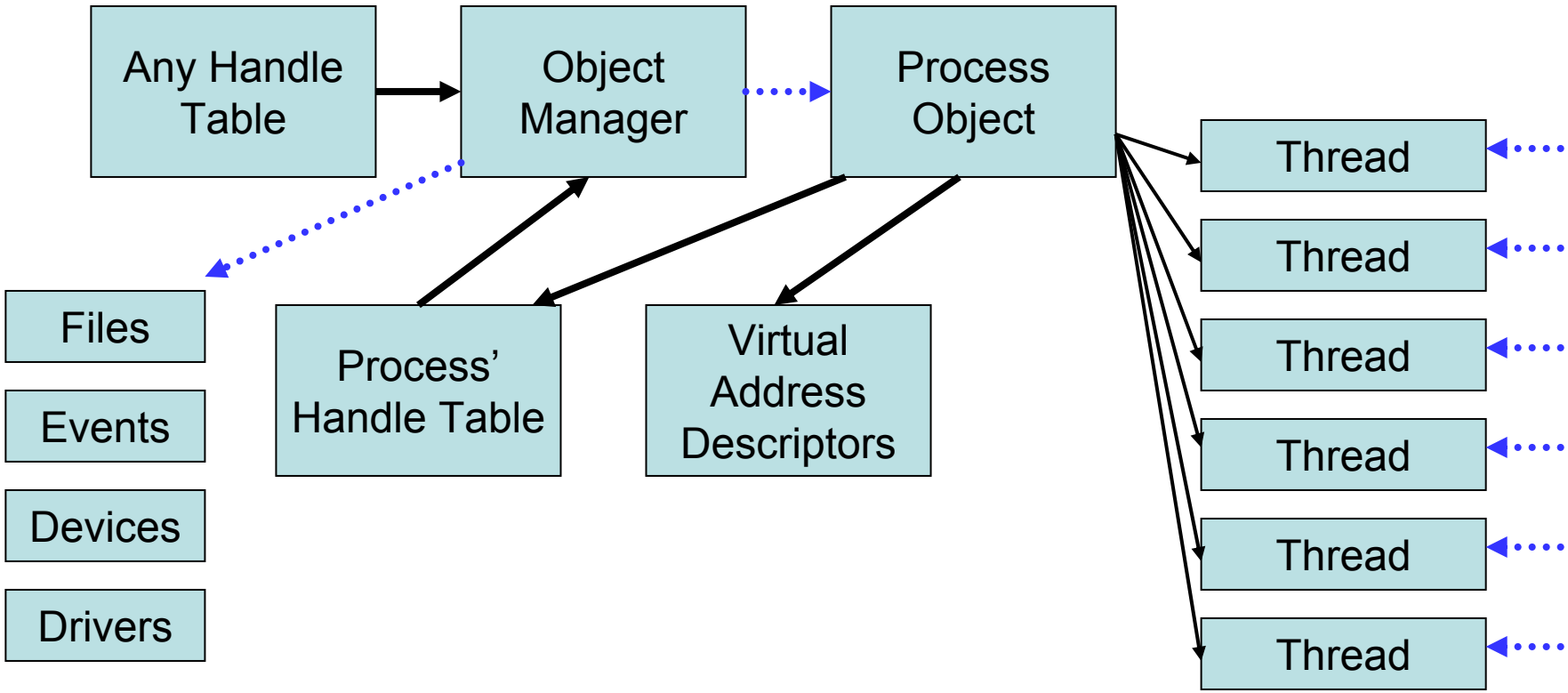
Scheduler (aka 'kernel')

Schedules thread execution on each processor

CPU Control-flow

- Thread scheduling occurs at PASSIVE or APC level (IRQL < 2)
- APCs (Asynchronous Procedure Calls) deliver I/O completions, thread/process termination, etc (IRQL == 1)
  - Not a general mechanism like unix signals (user-mode code must explicitly block pending APC delivery)
- Interrupt Service Routines run at IRL > 2
- ISRs defer most processing to run at IRQL==2 (DISPATCH level) by queuing a DPC to their current processor
- A pool of *worker threads* available for kernel components to run in a normal thread context when user-mode thread is unavailable or inappropriate
- Normal thread scheduling is round-robin among priority levels, with priority adjustments (except for fixed priority real-time threads)

Process/Thread structure



Process

Container for an address space and threads

Associated User-mode Process Environment Block (PEB)

Primary Access Token

Quota, Debug port, Handle Table etc

Unique process ID

Queued to the Job, global process list and Session list

MM structures like the WorkingSet, VAD tree, AWE etc

Thread

Fundamental schedulable entity in the system

Represented by ETHREAD that includes a KTHREAD

Queued to the process (both E and K thread)

IRP list

Impersonation Access Token

Unique thread ID

Associated User-mode Thread Environment Block (TEB)

User-mode stack

Kernel-mode stack

Processor Control Block (in KTHREAD) for cpu state when not running

Windows Past, Present, Future

PAST: Personal computer, 16->32 bits, MSDOS, Windows 9x code base, desktop focus

- Features, usability, compatibility, platform
- Windows 98

PRESENT: Enterprise computing, 32/64 bits, NT code base, solid desktop, datacenter

- Reliability, performance, IT Features
- Windows XP, Windows Server 2003

FUTURE: Managed code **(.NET Framework)**

- Productivity, innovation, empowerment
- Longhorn

.Net: Making it Simple

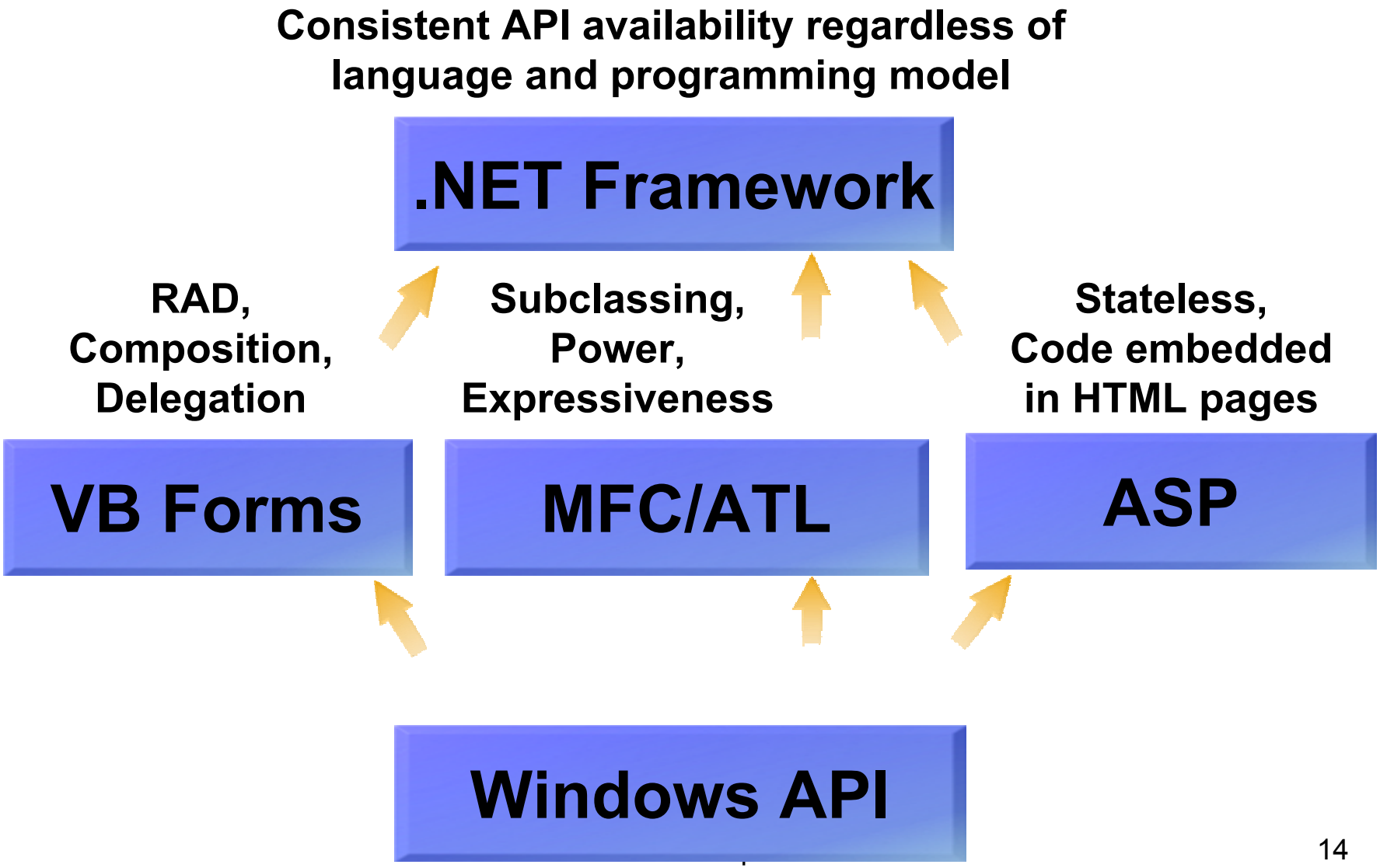
Windows API

```
HWND hwndMain = CreateWindowEx(
    0, "MainWClass", "Main Window",
    WS_OVERLAPPEDWINDOW | WS_HSCROLL | WS_VSCROLL,
    CW_USEDEFAULT, CW_USEDEFAULT,
    CW_USEDEFAULT, CW_USEDEFAULT,
    (HWND)NULL, (HMENU)NULL, hInstance, NULL);
ShowWindow(hwndMain, SW_SHOWDEFAULT);
UpdateWindow(hwndMain);
```

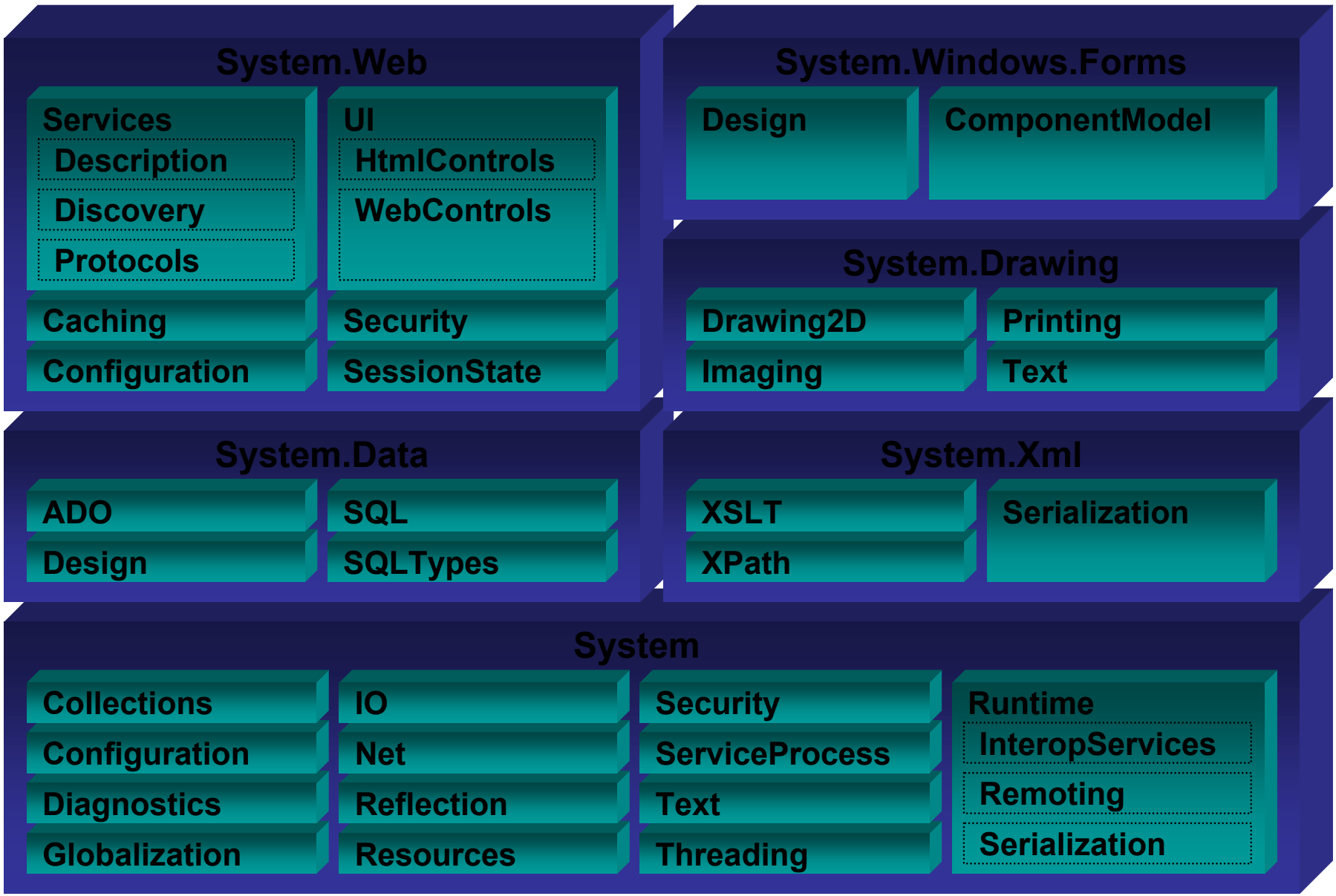
.Net Framework

```
Window w = new Window();
w.Text = "Main Window";
w.Show();
```

.Net: Unify Programming Models



.Net: API Organization



.Net: Languages

- The Managed Platform is Language Neutral
  - All languages are first class players
  - You can leverage your existing skills
- Common Language Specification
  - Set of features guaranteed to be in all languages
  - C# enforcement: [assembly:CLSCompliant(true)]
- We are providing
  - VB, C++, C#, J#, JScript
- Third-parties are building
  - APL, COBOL, Pascal, Eiffel, Haskell, ML, Oberon, Perl, Python, Scheme, Smalltalk…

Unmanaged vs. Managed

Unmanaged Code	Managed Code
Binary standard	**Type standard**
Type libraries	**Assemblies**
Immutable	**Resilient bind**
Reference counting	**Garbage collection**
Type unsafe	**Type safe**
Interface based	**Object based**
HRESULTs	**Exceptions**
GUIDs	**Strong names**

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Lectures

- Object Manager
- Virtual Memory
- Thread Scheduling
- Synchronization
- I/O Manager
- I/O Security
- Power Management
- NT File System
- Registry
- Lightweight Proc Calls
- Windows Services
- System Bootstrap
- Traps / Ints / Exceptions
- Processes
- Adv. Virtual Memory
- Cache Manager
- User-mode heap
- Win32k.sys
- WoW64
- Common Errors

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Projects

Device Drivers and Registry Hooking

Dragos Sambotin – Polytech. Inst. of Bucharest

Using LPC to build native client/server apps

Adrian Marinescu – University of Bucharest

Threads and Fibers

Arun Kishan – Stanford University

Doing virtual memory experiments from user-mode

Arun Kishan – Stanford University

Discussion